

# Chapter 17. Two Point Boundary Value Problems

## 17.0 Introduction

When ordinary differential equations are required to satisfy boundary conditions at more than one value of the independent variable, the resulting problem is called a *two point boundary value problem*. As the terminology indicates, the most common case by far is where boundary conditions are supposed to be satisfied at two points — usually the starting and ending values of the integration. However, the phrase “two point boundary value problem” is also used loosely to include more complicated cases, e.g., where some conditions are specified at endpoints, others at interior (usually singular) points.

The crucial distinction between initial value problems (Chapter 16) and two point boundary value problems (this chapter) is that in the former case we are able to start an acceptable solution at its beginning (initial values) and just march it along by numerical integration to its end (final values); while in the present case, the boundary conditions at the starting point do not determine a unique solution to start with — and a “random” choice among the solutions that satisfy these (incomplete) starting boundary conditions is almost certain *not* to satisfy the boundary conditions at the other specified point(s).

It should not surprise you that iteration is in general required to meld these spatially scattered boundary conditions into a single global solution of the differential equations. For this reason, two point boundary value problems require considerably more effort to solve than do initial value problems. You have to integrate your differential equations over the interval of interest, or perform an analogous “relaxation” procedure (see below), at least several, and sometimes very many, times. Only in the special case of linear differential equations can you say in advance just how many such iterations will be required.

The “standard” two point boundary value problem has the following form: We desire the solution to a set of  $N$  coupled first-order ordinary differential equations, satisfying  $n_1$  boundary conditions at the starting point  $x_1$ , and a remaining set of  $n_2 = N - n_1$  boundary conditions at the final point  $x_2$ . (Recall that all differential equations of order higher than first can be written as coupled sets of first-order equations, cf. §16.0.)

The differential equations are

$$\frac{dy_i(x)}{dx} = g_i(x, y_1, y_2, \dots, y_N) \quad i = 1, 2, \dots, N \quad (17.0.1)$$

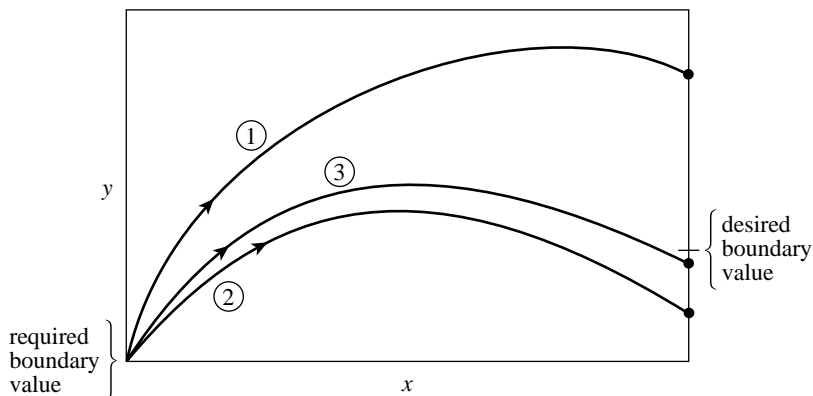


Figure 17.0.1. Shooting method (schematic). Trial integrations that satisfy the boundary condition at one endpoint are “launched.” The discrepancies from the desired boundary condition at the other endpoint are used to adjust the starting conditions, until boundary conditions at both endpoints are ultimately satisfied.

At  $x_1$ , the solution is supposed to satisfy

$$B_{1j}(x_1, y_1, y_2, \dots, y_N) = 0 \quad j = 1, \dots, n_1 \quad (17.0.2)$$

while at  $x_2$ , it is supposed to satisfy

$$B_{2k}(x_2, y_1, y_2, \dots, y_N) = 0 \quad k = 1, \dots, n_2 \quad (17.0.3)$$

There are two distinct classes of numerical methods for solving two point boundary value problems. In the *shooting method* (§17.1) we choose values for all of the dependent variables at one boundary. These values must be consistent with any boundary conditions for *that* boundary, but otherwise are arranged to depend on arbitrary free parameters whose values we initially “randomly” guess. We then integrate the ODEs by initial value methods, arriving at the other boundary (and/or any interior points with boundary conditions specified). In general, we find discrepancies from the desired boundary values there. Now we have a multidimensional root-finding problem, as was treated in §9.6 and §9.7: Find the adjustment of the free parameters at the starting point that zeros the discrepancies at the other boundary point(s). If we liken integrating the differential equations to following the trajectory of a shot from gun to target, then picking the initial conditions corresponds to aiming (see Figure 17.0.1). The shooting method provides a systematic approach to taking a set of “ranging” shots that allow us to improve our “aim” systematically.

As another variant of the shooting method (§17.2), we can guess unknown free parameters at both ends of the domain, integrate the equations to a common midpoint, and seek to adjust the guessed parameters so that the solution joins “smoothly” at the fitting point. In all shooting methods, trial solutions satisfy the differential equations “exactly” (or as exactly as we care to make our numerical integration), but the trial solutions come to satisfy the required boundary conditions only after the iterations are finished.

*Relaxation methods* use a different approach. The differential equations are replaced by finite-difference equations on a mesh of points that covers the range of

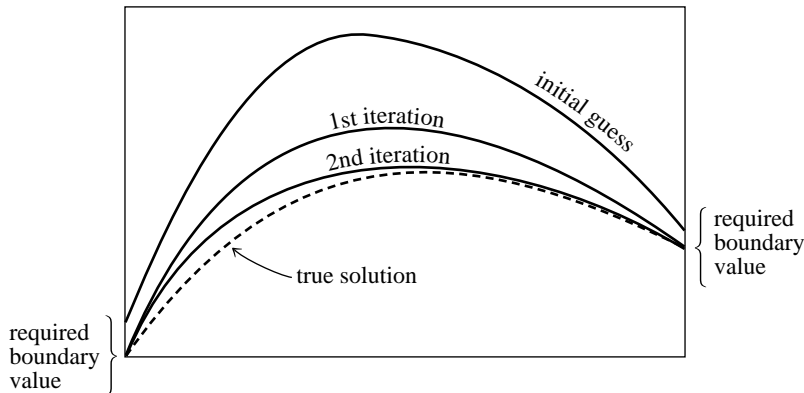


Figure 17.0.2. Relaxation method (schematic). An initial solution is guessed that approximately satisfies the differential equation and boundary conditions. An iterative process adjusts the function to bring it into close agreement with the true solution.

the integration. A trial solution consists of values for the dependent variables at each mesh point, *not* satisfying the desired finite-difference equations, nor necessarily even satisfying the required boundary conditions. The iteration, now called *relaxation*, consists of adjusting all the values on the mesh so as to bring them into successively closer agreement with the finite-difference equations and, simultaneously, with the boundary conditions (see Figure 17.0.2). For example, if the problem involves three coupled equations and a mesh of one hundred points, we must guess and improve three hundred variables representing the solution.

With all this adjustment, you may be surprised that relaxation is ever an efficient method, but (for the right problems) it really is! Relaxation works better than shooting when the boundary conditions are especially delicate or subtle, or where they involve complicated algebraic relations that cannot easily be solved in closed form. Relaxation works best when the solution is smooth and not highly oscillatory. Such oscillations would require many grid points for accurate representation. The number and position of required points may not be known *a priori*. Shooting methods are usually preferred in such cases, because their variable stepsize integrations adjust naturally to a solution's peculiarities.

Relaxation methods are often preferred when the ODEs have extraneous solutions which, while not appearing in the final solution satisfying all boundary conditions, may wreak havoc on the initial value integrations required by shooting. The typical case is that of trying to maintain a dying exponential in the presence of growing exponentials.

Good initial guesses are the secret of efficient relaxation methods. Often one has to solve a problem many times, each time with a slightly different value of some parameter. In that case, the previous solution is usually a good initial guess when the parameter is changed, and relaxation will work well.

Until you have enough experience to make your own judgment between the two methods, you might wish to follow the advice of your authors, who are notorious computer gunslingers: We always shoot first, and only then relax.

### Problems Reducible to the Standard Boundary Problem

There are two important problems that can be reduced to the standard boundary value problem described by equations (17.0.1) – (17.0.3). The first is the *eigenvalue problem for differential equations*. Here the right-hand side of the system of differential equations depends on a parameter  $\lambda$ ,

$$\frac{dy_i(x)}{dx} = g_i(x, y_1, \dots, y_N, \lambda) \quad (17.0.4)$$

and one has to satisfy  $N + 1$  boundary conditions instead of just  $N$ . The problem is overdetermined and in general there is no solution for arbitrary values of  $\lambda$ . For certain special values of  $\lambda$ , the eigenvalues, equation (17.0.4) does have a solution.

We reduce this problem to the standard case by introducing a new dependent variable

$$y_{N+1} \equiv \lambda \quad (17.0.5)$$

and another differential equation

$$\frac{dy_{N+1}}{dx} = 0 \quad (17.0.6)$$

An example of this trick is given in §17.4.

The other case that can be put in the standard form is a *free boundary problem*. Here only one boundary abscissa  $x_1$  is specified, while the other boundary  $x_2$  is to be determined so that the system (17.0.1) has a solution satisfying a total of  $N + 1$  boundary conditions. Here we again add an extra constant dependent variable:

$$y_{N+1} \equiv x_2 - x_1 \quad (17.0.7)$$

$$\frac{dy_{N+1}}{dx} = 0 \quad (17.0.8)$$

We also define a new *independent* variable  $t$  by setting

$$x - x_1 \equiv t y_{N+1}, \quad 0 \leq t \leq 1 \quad (17.0.9)$$

The system of  $N + 1$  differential equations for  $dy_i/dt$  is now in the standard form, with  $t$  varying between the known limits 0 and 1.

#### CITED REFERENCES AND FURTHER READING:

- Keller, H.B. 1968, *Numerical Methods for Two-Point Boundary-Value Problems* (Waltham, MA: Blaisdell).
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- Eggleton, P.P. 1971, *Monthly Notices of the Royal Astronomical Society*, vol. 151, pp. 351–364.
- London, R.A., and Flannery, B.P. 1982, *Astrophysical Journal*, vol. 258, pp. 260–269.
- Stoer, J., and Bulirsch, R. 1980, *Introduction to Numerical Analysis* (New York: Springer-Verlag), §§7.3–7.4.

## 17.1 The Shooting Method

In this section we discuss “pure” shooting, where the integration proceeds from  $x_1$  to  $x_2$ , and we try to match boundary conditions at the end of the integration. In the next section, we describe shooting to an intermediate fitting point, where the solution to the equations and boundary conditions is found by launching “shots” from both sides of the interval and trying to match continuity conditions at some intermediate point.

Our implementation of the shooting method exactly implements multidimensional, globally convergent Newton-Raphson (§9.7). It seeks to zero  $n_2$  functions of  $n_2$  variables. The functions are obtained by integrating  $N$  differential equations from  $x_1$  to  $x_2$ . Let us see how this works:

At the starting point  $x_1$  there are  $N$  starting values  $y_i$  to be specified, but subject to  $n_1$  conditions. Therefore there are  $n_2 = N - n_1$  *freely specifiable* starting values. Let us imagine that these freely specifiable values are the components of a vector  $\mathbf{V}$  that lives in a vector space of dimension  $n_2$ . Then you, the user, knowing the functional form of the boundary conditions (17.0.2), can write a subroutine that generates a complete set of  $N$  starting values  $\mathbf{y}$ , satisfying the boundary conditions at  $x_1$ , from an arbitrary vector value of  $\mathbf{V}$  in which there are no restrictions on the  $n_2$  component values. In other words, (17.0.2) converts to a prescription

$$y_i(x_1) = y_i(x_1; V_1, \dots, V_{n_2}) \quad i = 1, \dots, N \quad (17.1.1)$$

Below, the subroutine that implements (17.1.1) will be called `load`.

Notice that the components of  $\mathbf{V}$  might be exactly the values of certain “free” components of  $\mathbf{y}$ , with the other components of  $\mathbf{y}$  determined by the boundary conditions. Alternatively, the components of  $\mathbf{V}$  might parametrize the solutions that satisfy the starting boundary conditions in some other convenient way. Boundary conditions often impose algebraic relations among the  $y_i$ , rather than specific values for each of them. Using some auxiliary set of parameters often makes it easier to “solve” the boundary relations for a consistent set of  $y_i$ ’s. It makes no difference which way you go, as long as your vector space of  $\mathbf{V}$ ’s generates (through 17.1.1) all allowed starting vectors  $\mathbf{y}$ .

Given a particular  $\mathbf{V}$ , a particular  $\mathbf{y}(x_1)$  is thus generated. It can then be turned into a  $\mathbf{y}(x_2)$  by integrating the ODEs to  $x_2$  as an initial value problem (e.g., using Chapter 16’s `odeint`). Now, at  $x_2$ , let us define a *discrepancy vector*  $\mathbf{F}$ , also of dimension  $n_2$ , whose components measure how far we are from satisfying the  $n_2$  boundary conditions at  $x_2$  (17.0.3). Simplest of all is just to use the right-hand sides of (17.0.3),

$$F_k = B_{2k}(x_2, \mathbf{y}) \quad k = 1, \dots, n_2 \quad (17.1.2)$$

As in the case of  $\mathbf{V}$ , however, you can use any other convenient parametrization, as long as your space of  $\mathbf{F}$ ’s spans the space of possible discrepancies from the desired boundary conditions, with all components of  $\mathbf{F}$  equal to zero if and only if the boundary conditions at  $x_2$  are satisfied. Below, you will be asked to supply a user-written subroutine `score` which uses (17.0.3) to convert an  $N$ -vector of ending values  $\mathbf{y}(x_2)$  into an  $n_2$ -vector of discrepancies  $\mathbf{F}$ .

Now, as far as Newton-Raphson is concerned, we are nearly in business. We want to find a vector value of  $\mathbf{V}$  that zeros the vector value of  $\mathbf{F}$ . We do this by invoking the globally convergent Newton's method implemented in the routine `newt` of §9.7. Recall that the heart of Newton's method involves solving the set of  $n_2$  linear equations

$$\mathbf{J} \cdot \delta\mathbf{V} = -\mathbf{F} \quad (17.1.3)$$

and then adding the correction back,

$$\mathbf{V}^{\text{new}} = \mathbf{V}^{\text{old}} + \delta\mathbf{V} \quad (17.1.4)$$

In (17.1.3), the Jacobian matrix  $\mathbf{J}$  has components given by

$$J_{ij} = \frac{\partial F_i}{\partial V_j} \quad (17.1.5)$$

It is not feasible to compute these partial derivatives analytically. Rather, each requires a *separate* integration of the  $N$  ODEs, followed by the evaluation of

$$\frac{\partial F_i}{\partial V_j} \approx \frac{F_i(V_1, \dots, V_j + \Delta V_j, \dots) - F_i(V_1, \dots, V_j, \dots)}{\Delta V_j} \quad (17.1.6)$$

This is done automatically for you in the routine `fdjac` that comes with `newt`. The only input to `newt` that you have to provide is the routine `funcv` that calculates  $\mathbf{F}$  by integrating the ODEs. Here is the appropriate routine:

```

C SUBROUTINE shoot(n2,v,f) is named "funcv" for use with "newt"
SUBROUTINE funcv(n2,v,f)
  INTEGER n2,nvar,kmax,kount,KMAXX,NMAX
  REAL f(n2),v(n2),x1,x2,dxsav,yp,eps
  PARAMETER (NMAX=50,KMAXX=200,EPS=1.e-6) At most NMAX coupled ODEs.
  COMMON /caller/ x1,x2,nvar
  COMMON /path/ kmax,kount,dxsav,yp(KMAXX),yp(NMAX,KMAXX)
C USES derivs,load,odeint,rkqs,score
  Routine for use with newt to solve a two point boundary value problem for nvar coupled
  ODEs by shooting from x1 to x2. Initial values for the nvar ODEs at x1 are generated
  from the n2 input coefficients v(1:n2), using the user-supplied routine load. The routine
  integrates the ODEs to x2 using the Runge-Kutta method with tolerance EPS, initial stepsize
  h1, and minimum stepsize hmin. At x2 it calls the user-supplied subroutine score to
  evaluate the n2 functions f(1:n2) that ought to be zero to satisfy the boundary conditions
  at x2. The functions f are returned on output. newt uses a globally convergent Newton's
  method to adjust the values of v until the functions f are zero. The user-supplied subroutine
  derivs(x,y,dydx) supplies derivative information to the ODE integrator (see Chapter
  16). The common block caller receives its values from the main program so that funcv
  can have the syntax required by newt. The common block path is included for compatibility
  with odeint.
  INTEGER nbad,nok
  REAL h1,hmin,y(NMAX)
  EXTERNAL derivs,rkqs
  kmax=0
  h1=(x2-x1)/100.
  hmin=0.
  call load(x1,v,y)
  call odeint(y,nvar,x1,x2,eps,h1,hmin,nok,nbad,derivs,rkqs)
  call score(x2,y,f)
  return
END

```

For some problems the initial stepsize  $\Delta V$  might depend sensitively upon the initial conditions. It is straightforward to alter `load` to include a suggested stepsize `h1` as another returned argument and feed it to `fdjac` via a common block.

A complete cycle of the shooting method thus requires  $n_2 + 1$  integrations of the  $N$  coupled ODEs: one integration to evaluate the current degree of mismatch, and  $n_2$  for the partial derivatives. Each new cycle requires a new round of  $n_2 + 1$  integrations. This illustrates the enormous extra effort involved in solving two point boundary value problems compared with initial value problems.

If the differential equations are *linear*, then only one complete cycle is required, since (17.1.3)–(17.1.4) should take us right to the solution. A second round can be useful, however, in mopping up some (never all) of the roundoff error.

As given here, `shoot` uses the quality controlled Runge-Kutta method of §16.2 to integrate the ODEs, but any of the other methods of Chapter 16 could just as well be used.

You, the user, must supply `shoot` with: (i) a subroutine `load(x1, v, y)` which returns the  $n$ -vector  $y(1:n)$  (satisfying the starting boundary conditions, of course), given the freely specifiable variables of  $v(1:n2)$  at the initial point  $x1$ ; (ii) a subroutine `score(x2, y, f)` which returns the discrepancy vector  $f(1:n2)$  of the ending boundary conditions, given the vector  $y(1:n)$  at the endpoint  $x2$ ; (iii) a starting vector  $v(1:n2)$ ; (iv) a subroutine `derivs` for the ODE integration; and other obvious parameters as described in the header comment above.

In §17.4 we give a sample program illustrating how to use `shoot`.

#### CITED REFERENCES AND FURTHER READING:

- Acton, F.S. 1970, *Numerical Methods That Work*; 1990, corrected edition (Washington: Mathematical Association of America).  
 Keller, H.B. 1968, *Numerical Methods for Two-Point Boundary-Value Problems* (Waltham, MA: Blaisdell).

## 17.2 Shooting to a Fitting Point

The shooting method described in §17.1 tacitly assumed that the “shots” would be able to traverse the entire domain of integration, even at the early stages of convergence to a correct solution. In some problems it can happen that, for very wrong starting conditions, an initial solution can’t even get from  $x_1$  to  $x_2$  without encountering some incalculable, or catastrophic, result. For example, the argument of a square root might go negative, causing the numerical code to crash. Simple shooting would be stymied.

A different, but related, case is where the endpoints are both singular points of the set of ODEs. One frequently needs to use special methods to integrate near the singular points, analytic asymptotic expansions, for example. In such cases it is feasible to integrate in the direction *away* from a singular point, using the special method to get through the first little bit and then reading off “initial” values for further numerical integration. However it is usually not feasible to integrate *into* a singular point, if only because one has not usually expended the same analytic

effort to obtain expansions of “wrong” solutions near the singular point (those not satisfying the desired boundary condition).

The solution to the above mentioned difficulties is *shooting to a fitting point*. Instead of integrating from  $x_1$  to  $x_2$ , we integrate first from  $x_1$  to some point  $x_f$  that is *between*  $x_1$  and  $x_2$ ; and second from  $x_2$  (in the opposite direction) to  $x_f$ .

If (as before) the number of boundary conditions imposed at  $x_1$  is  $n_1$ , and the number imposed at  $x_2$  is  $n_2$ , then there are  $n_2$  freely specifiable starting values at  $x_1$  and  $n_1$  freely specifiable starting values at  $x_2$ . (If you are confused by this, go back to §17.1.) We can therefore define an  $n_2$ -vector  $\mathbf{V}_{(1)}$  of starting parameters at  $x_1$ , and a prescription `load1(x1, v1, y)` for mapping  $\mathbf{V}_{(1)}$  into a  $\mathbf{y}$  that satisfies the boundary conditions at  $x_1$ ,

$$y_i(x_1) = y_i(x_1; V_{(1)1}, \dots, V_{(1)n_2}) \quad i = 1, \dots, N \quad (17.2.1)$$

Likewise we can define an  $n_1$ -vector  $\mathbf{V}_{(2)}$  of starting parameters at  $x_2$ , and a prescription `load2(x2, v2, y)` for mapping  $\mathbf{V}_{(2)}$  into a  $\mathbf{y}$  that satisfies the boundary conditions at  $x_2$ ,

$$y_i(x_2) = y_i(x_2; V_{(2)1}, \dots, V_{(2)n_1}) \quad i = 1, \dots, N \quad (17.2.2)$$

We thus have a total of  $N$  freely adjustable parameters in the combination of  $\mathbf{V}_{(1)}$  and  $\mathbf{V}_{(2)}$ . The  $N$  conditions that must be satisfied are that there be agreement in  $N$  components of  $\mathbf{y}$  at  $x_f$  between the values obtained integrating from one side and from the other,

$$y_i(x_f; \mathbf{V}_{(1)}) = y_i(x_f; \mathbf{V}_{(2)}) \quad i = 1, \dots, N \quad (17.2.3)$$

In some problems, the  $N$  matching conditions can be better described (physically, mathematically, or numerically) by using  $N$  different functions  $F_i$ ,  $i = 1 \dots N$ , each possibly depending on the  $N$  components  $y_i$ . In those cases, (17.2.3) is replaced by

$$F_i[\mathbf{y}(x_f; \mathbf{V}_{(1)})] = F_i[\mathbf{y}(x_f; \mathbf{V}_{(2)})] \quad i = 1, \dots, N \quad (17.2.4)$$

In the program below, the user-supplied subroutine `score(xf, y, f)` is supposed to map an input  $N$ -vector  $\mathbf{y}$  into an output  $N$ -vector  $\mathbf{F}$ . In most cases, you can dummy this subroutine as the identity mapping.

Shooting to a fitting point uses globally convergent Newton-Raphson exactly as in §17.1. Comparing closely with the routine `shoot` of the previous section, you should have no difficulty in understanding the following routine `shootf`. The main differences in use are that you have to supply both `load1` and `load2`. Also, in the calling program you must supply initial guesses for `v1(1:n2)` and `v2(1:n1)`. Once again a sample program illustrating shooting to a fitting point is given in §17.4.

```

C  SUBROUTINE shootf(n,v,f) is named "funcv" for use with "newt"
SUBROUTINE funcv(n,v,f)
INTEGER n,nvar,nn2,kmax,kount,KMAXX,NMAX
REAL f(n),v(n),x1,x2,xf,dxsav,yp,eps
PARAMETER (NMAX=50,KMAXX=200,EPS=1.e-6) At most NMAX equations.
COMMON /caller/ x1,x2,xf,nvar,nn2
COMMON /path/ kmax,kount,dxsav,yp(KMAXX),yp(NMAX,KMAXX)
C  USES derivs,load1,load2,odeint,rkqs,score
Routine for use with newt to solve a two point boundary value problem for nvar coupled ODEs by shooting from x1 and x2 to a fitting point xf. Initial values for the nvar

```



ODEs at  $x_1$  ( $x_2$ ) are generated from the  $n_2$  ( $n_1$ ) coefficients  $v_1$  ( $v_2$ ), using the user-supplied routine `load1` (`load2`). The coefficients  $v_1$  and  $v_2$  should be stored in a single array  $v(1:n_1+n_2)$  in the main program by an EQUIVALENCE statement of the form  $(v_1(1), v(1)), (v_2(1), v(n_2+1))$ . The input parameter  $n = n_1 + n_2 = nvar$ . The routine integrates the ODEs to  $xf$  using the Runge-Kutta method with tolerance `EPS`, initial stepsize `h1`, and minimum stepsize `hmin`. At  $xf$  it calls the user-supplied subroutine `score` to evaluate the  $nvar$  functions `f1` and `f2` that ought to match at  $xf$ . The differences  $f$  are returned on output. `newt` uses a globally convergent Newton's method to adjust the values of  $v$  until the functions  $f$  are zero. The user-supplied subroutine `derivs(x,y,dydx)` supplies derivative information to the ODE integrator (see Chapter 16). The common block `caller` receives its values from the main program so that `funcv` can have the syntax required by `newt`. Set `nn2 = n2` in the main program. The common block path is for compatibility with `odeint`.

```

INTEGER i,nbad,nok
REAL h1,hmin,f1(NMAX),f2(NMAX),y(NMAX)
EXTERNAL derivs,rkqs
kmax=0
h1=(x2-x1)/100.
hmin=0.
call load1(x1,v,y)           Path from x1 to xf with best trial values v1.
call odeint(y,nvar,x1,xf,EPS,h1,hmin,nok,nbad,derivs,rkqs)
call score(xf,y,f1)
call load2(x2,v(nn2+1),y)    Path from x2 to xf with best trial values v2.
call odeint(y,nvar,x2,xf,EPS,h1,hmin,nok,nbad,derivs,rkqs)
call score(xf,y,f2)
do 11 i=1,n
  f(i)=f1(i)-f2(i)
enddo 11
return
END

```

There are boundary value problems where even shooting to a fitting point fails — the integration interval has to be partitioned by several fitting points with the solution being matched at each such point. For more details see [1].

#### CITED REFERENCES AND FURTHER READING:

- Acton, F.S. 1970, *Numerical Methods That Work*; 1990, corrected edition (Washington: Mathematical Association of America).
- Keller, H.B. 1968, *Numerical Methods for Two-Point Boundary-Value Problems* (Waltham, MA: Blaisdell).
- Stoer, J., and Bulirsch, R. 1980, *Introduction to Numerical Analysis* (New York: Springer-Verlag), §§7.3.5–7.3.6. [1]

## 17.3 Relaxation Methods

In *relaxation methods* we replace ODEs by approximate *finite-difference equations* (FDEs) on a grid or mesh of points that spans the domain of interest. As a typical example, we could replace a general first-order differential equation

$$\frac{dy}{dx} = g(x, y) \quad (17.3.1)$$

with an algebraic equation relating function values at two points  $k, k-1$ :

$$y_k - y_{k-1} - (x_k - x_{k-1}) g \left[ \frac{1}{2}(x_k + x_{k-1}), \frac{1}{2}(y_k + y_{k-1}) \right] = 0 \quad (17.3.2)$$

The form of the FDE in (17.3.2) illustrates the idea, but not uniquely: There are many ways to turn the ODE into an FDE. When the problem involves  $N$  coupled first-order ODEs represented by FDEs on a mesh of  $M$  points, a solution consists of values for  $N$  dependent functions given at each of the  $M$  mesh points, or  $N \times M$  variables in all. The relaxation method determines the solution by starting with a guess and improving it, iteratively. As the iterations improve the solution, the result is said to *relax* to the true solution.

While several iteration schemes are possible, for most problems our old standby, multi-dimensional Newton's method, works well. The method produces a matrix equation that must be solved, but the matrix takes a special, "block diagonal" form, that allows it to be inverted far more economically both in time and storage than would be possible for a general matrix of size  $(MN) \times (MN)$ . Since  $MN$  can easily be several thousand, this is crucial for the feasibility of the method.

Our implementation couples at most pairs of points, as in equation (17.3.2). More points can be coupled, but then the method becomes more complex. We will provide enough background so that you can write a more general scheme if you have the patience to do so.

Let us develop a general set of algebraic equations that represent the ODEs by FDEs. The ODE problem is exactly identical to that expressed in equations (17.0.1)–(17.0.3) where we had  $N$  coupled first-order equations that satisfy  $n_1$  boundary conditions at  $x_1$  and  $n_2 = N - n_1$  boundary conditions at  $x_2$ . We first define a mesh or grid by a set of  $k = 1, 2, \dots, M$  points at which we supply values for the independent variable  $x_k$ . In particular,  $x_1$  is the initial boundary, and  $x_M$  is the final boundary. We use the notation  $\mathbf{y}_k$  to refer to the entire set of dependent variables  $y_1, y_2, \dots, y_N$  at point  $x_k$ . At an arbitrary point  $k$  in the middle of the mesh, we approximate the set of  $N$  first-order ODEs by algebraic relations of the form

$$0 = \mathbf{E}_k \equiv \mathbf{y}_k - \mathbf{y}_{k-1} - (x_k - x_{k-1})\mathbf{g}_k(x_k, x_{k-1}, \mathbf{y}_k, \mathbf{y}_{k-1}), \quad k = 2, 3, \dots, M \quad (17.3.3)$$

The notation signifies that  $\mathbf{g}_k$  can be evaluated using information from both points  $k, k - 1$ . The FDEs labeled by  $\mathbf{E}_k$  provide  $N$  equations coupling  $2N$  variables at points  $k, k - 1$ . There are  $M - 1$  points,  $k = 2, 3, \dots, M$ , at which difference equations of the form (17.3.3) apply. Thus the FDEs provide a total of  $(M - 1)N$  equations for the  $MN$  unknowns. The remaining  $N$  equations come from the boundary conditions.

At the first boundary we have

$$0 = \mathbf{E}_1 \equiv \mathbf{B}(x_1, \mathbf{y}_1) \quad (17.3.4)$$

while at the second boundary

$$0 = \mathbf{E}_{M+1} \equiv \mathbf{C}(x_M, \mathbf{y}_M) \quad (17.3.5)$$

The vectors  $\mathbf{E}_1$  and  $\mathbf{B}$  have only  $n_1$  nonzero components, corresponding to the  $n_1$  boundary conditions at  $x_1$ . It will turn out to be useful to take these nonzero components to be the *last*  $n_1$  components. In other words,  $E_{j,1} \neq 0$  only for  $j = n_2 + 1, n_2 + 2, \dots, N$ . At the other boundary, only the first  $n_2$  components of  $\mathbf{E}_{M+1}$  and  $\mathbf{C}$  are nonzero:  $E_{j,M+1} \neq 0$  only for  $j = 1, 2, \dots, n_2$ .

The "solution" of the FDE problem in (17.3.3)–(17.3.5) consists of a set of variables  $y_{j,k}$ , the values of the  $N$  variables  $y_j$  at the  $M$  points  $x_k$ . The algorithm we describe below requires an initial guess for the  $y_{j,k}$ . We then determine increments  $\Delta y_{j,k}$  such that  $y_{j,k} + \Delta y_{j,k}$  is an improved approximation to the solution.

Equations for the increments are developed by expanding the FDEs in first-order Taylor series with respect to small changes  $\Delta \mathbf{y}_k$ . At an interior point,  $k = 2, 3, \dots, M$  this gives:

$$\begin{aligned} \mathbf{E}_k(\mathbf{y}_k + \Delta \mathbf{y}_k, \mathbf{y}_{k-1} + \Delta \mathbf{y}_{k-1}) &\approx \mathbf{E}_k(\mathbf{y}_k, \mathbf{y}_{k-1}) \\ &+ \sum_{n=1}^N \frac{\partial \mathbf{E}_k}{\partial y_{n,k-1}} \Delta y_{n,k-1} + \sum_{n=1}^N \frac{\partial \mathbf{E}_k}{\partial y_{n,k}} \Delta y_{n,k} \end{aligned} \quad (17.3.6)$$

For a solution we want the updated value  $\mathbf{E}(\mathbf{y} + \Delta \mathbf{y})$  to be zero, so the general set of equations at an interior point can be written in matrix form as

$$\sum_{n=1}^N S_{j,n} \Delta y_{n,k-1} + \sum_{n=N+1}^{2N} S_{j,n} \Delta y_{n-N,k} = -E_{j,k}, \quad j = 1, 2, \dots, N \quad (17.3.7)$$

where

$$S_{j,n} = \frac{\partial E_{j,k}}{\partial y_{n,k-1}}, \quad S_{j,n+N} = \frac{\partial E_{j,k}}{\partial y_{n,k}}, \quad n = 1, 2, \dots, N \quad (17.3.8)$$

The quantity  $S_{j,n}$  is an  $N \times 2N$  matrix at each point  $k$ . Each interior point thus supplies a block of  $N$  equations coupling  $2N$  corrections to the solution variables at the points  $k, k-1$ .

Similarly, the algebraic relations at the boundaries can be expanded in a first-order Taylor series for increments that improve the solution. Since  $E_1$  depends only on  $y_1$ , we find at the first boundary:

$$\sum_{n=1}^N S_{j,n} \Delta y_{n,1} = -E_{j,1}, \quad j = n_2 + 1, n_2 + 2, \dots, N \quad (17.3.9)$$

where

$$S_{j,n} = \frac{\partial E_{j,1}}{\partial y_{n,1}}, \quad n = 1, 2, \dots, N \quad (17.3.10)$$

At the second boundary,

$$\sum_{n=1}^N S_{j,n} \Delta y_{n,M} = -E_{j,M+1}, \quad j = 1, 2, \dots, n_2 \quad (17.3.11)$$

where

$$S_{j,n} = \frac{\partial E_{j,M+1}}{\partial y_{n,M}}, \quad n = 1, 2, \dots, N \quad (17.3.12)$$

We thus have in equations (17.3.7)–(17.3.12) a set of linear equations to be solved for the corrections  $\Delta y$ , iterating until the corrections are sufficiently small. The equations have a special structure, because each  $S_{j,n}$  couples only points  $k, k-1$ . Figure 17.3.1 illustrates the typical structure of the complete matrix equation for the case of 5 variables and 4 mesh points, with 3 boundary conditions at the first boundary and 2 at the second. The  $3 \times 5$  block of nonzero entries in the top left-hand corner of the matrix comes from the boundary condition  $S_{j,n}$  at point  $k=1$ . The next three  $5 \times 10$  blocks are the  $S_{j,n}$  at the interior points, coupling variables at mesh points (2,1), (3,2), and (4,3). Finally we have the block corresponding to the second boundary condition.

We can solve equations (17.3.7)–(17.3.12) for the increments  $\Delta y$  using a form of Gaussian elimination that exploits the special structure of the matrix to minimize the total number of operations, and that minimizes storage of matrix coefficients by packing the elements in a special blocked structure. (You might wish to review Chapter 2, especially §2.2, if you are unfamiliar with the steps involved in Gaussian elimination.) Recall that Gaussian elimination consists of manipulating the equations by elementary operations such as dividing rows of coefficients by a common factor to produce unity in diagonal elements, and adding appropriate multiples of other rows to produce zeros below the diagonal. Here we take advantage of the block structure by performing a bit more reduction than in pure Gaussian elimination, so that the storage of coefficients is minimized. Figure 17.3.2 shows the form that we wish to achieve by elimination, just prior to the backsubstitution step. Only a small subset of the reduced  $MN \times MN$  matrix elements needs to be stored as the elimination progresses. Once the matrix elements reach the stage in Figure 17.3.2, the solution follows quickly by a backsubstitution procedure.

Furthermore, the entire procedure, except the backsubstitution step, operates only on one block of the matrix at a time. The procedure contains four types of operations: (1) partial reduction to zero of certain elements of a block using results from a previous step, (2) elimination of the square structure of the remaining block elements such that the square section contains unity along the diagonal, and zero in off-diagonal elements, (3) storage of the remaining nonzero coefficients for use in later steps, and (4) backsubstitution. We illustrate the steps schematically by figures.

Consider the block of equations describing corrections available from the initial boundary conditions. We have  $n_1$  equations for  $N$  unknown corrections. We wish to transform the first

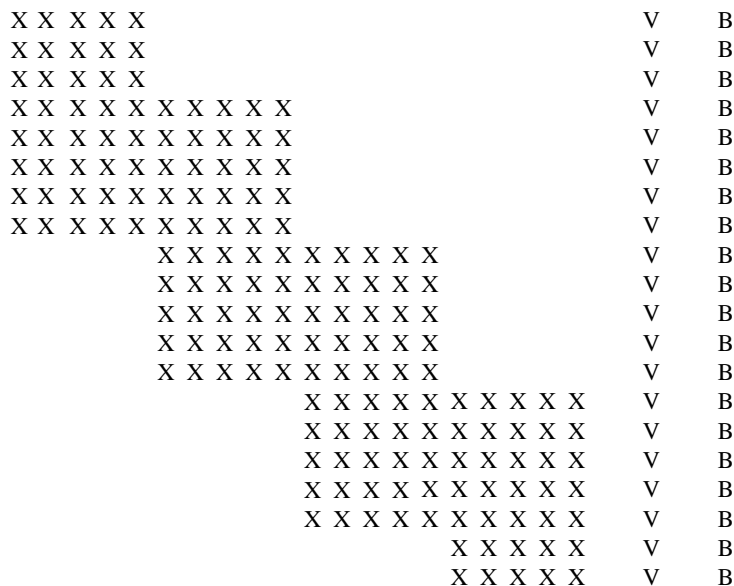


Figure 17.3.1. Matrix structure of a set of linear finite-difference equations (FDEs) with boundary conditions imposed at both endpoints. Here **X** represents a coefficient of the FDEs, **V** represents a component of the unknown solution vector, and **B** is a component of the known right-hand side. Empty spaces represent zeros. The matrix equation is to be solved by a special form of Gaussian elimination. (See text for details.)

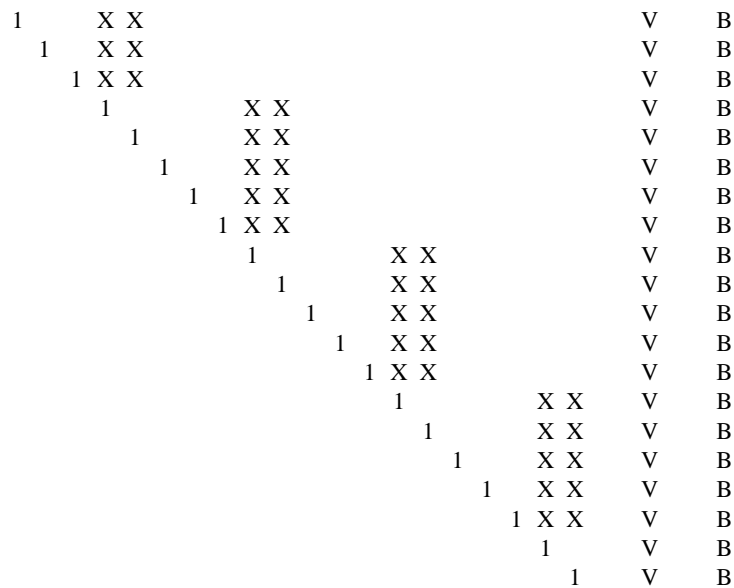


Figure 17.3.2. Target structure of the Gaussian elimination. Once the matrix of Figure 17.3.1 has been reduced to this form, the solution follows quickly by backsubstitution.

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block so that its left-hand  $n_1 \times n_1$  square section becomes unity along the diagonal, and zero in off-diagonal elements. Figure 17.3.3 shows the original and final form of the first block of the matrix. In the figure we designate matrix elements that are subject to diagonalization by “D”, and elements that will be altered by “A”; in the final block, elements that are stored are labeled by “S”. We get from start to finish by selecting in turn  $n_1$  “pivot” elements from among the first  $n_1$  columns, normalizing the pivot row so that the value of the “pivot” element is unity, and adding appropriate multiples of this row to the remaining rows so that they contain zeros in the pivot column. In its final form, the reduced block expresses values for the corrections to the first  $n_1$  variables at mesh point 1 in terms of values for the remaining  $n_2$  unknown corrections at point 1, i.e., we now know what the first  $n_1$  elements are in terms of the remaining  $n_2$  elements. We store only the final set of  $n_2$  nonzero columns from the initial block, plus the column for the altered right-hand side of the matrix equation.

We must emphasize here an important detail of the method. To exploit the reduced storage allowed by operating on blocks, it is essential that the ordering of columns in the  $\mathbf{s}$  matrix of derivatives be such that pivot elements can be found among the first  $n_1$  rows of the matrix. This means that the  $n_1$  boundary conditions at the first point must contain some dependence on the first  $j=1, 2, \dots, n_1$  dependent variables,  $y(j, 1)$ . If not, then the original square  $n_1 \times n_1$  subsection of the first block will appear to be singular, and the method will fail. Alternatively, we would have to allow the search for pivot elements to involve all  $N$  columns of the block, and this would require column swapping and far more bookkeeping. The code provides a simple method of reordering the variables, i.e., the columns of the  $\mathbf{s}$  matrix, so that this can be done easily. End of important detail.

Next consider the block of  $N$  equations representing the FDEs that describe the relation between the  $2N$  corrections at points 2 and 1. The elements of that block, together with results from the previous step, are illustrated in Figure 17.3.4. Note that by adding suitable multiples of rows from the first block we can reduce to zero the first  $n_1$  columns of the block (labeled by “Z”), and, to do so, we will need to alter only the columns from  $n_1 + 1$  to  $N$  and the vector element on the right-hand side. Of the remaining columns we can diagonalize a square subsection of  $N \times N$  elements, labeled by “D” in the figure. In the process we alter the final set of  $n_2 + 1$  columns, denoted “A” in the figure. The second half of the figure shows the block when we finish operating on it, with the stored  $(n_2 + 1) \times N$  elements labeled by “S.”

If we operate on the next set of equations corresponding to the FDEs coupling corrections at points 3 and 2, we see that the state of available results and new equations exactly reproduces the situation described in the previous paragraph. Thus, we can carry out those steps again for each block in turn through block  $M$ . Finally on block  $M + 1$  we encounter the remaining boundary conditions.

Figure 17.3.5 shows the final block of  $n_2$  FDEs relating the  $N$  corrections for variables at mesh point  $M$ , together with the result of reducing the previous block. Again, we can first use the prior results to zero the first  $n_1$  columns of the block. Now, when we diagonalize the remaining square section, we strike gold: We get values for the final  $n_2$  corrections at mesh point  $M$ .

With the final block reduced, the matrix has the desired form shown previously in Figure 17.3.2, and the matrix is ripe for backsubstitution. Starting with the bottom row and working up towards the top, at each stage we can simply determine one unknown correction in terms of known quantities.

The subroutine `solvde` organizes the steps described above. The principal procedures used in the algorithm are performed by subroutines called internally by `solvde`. The subroutine `red` eliminates leading columns of the  $\mathbf{s}$  matrix using results from prior blocks. `pinvs` diagonalizes the square subsection of  $\mathbf{s}$  and stores unreduced coefficients. `bksub` carries out the backsubstitution step. The user of `solvde` must understand the calling arguments, as described below, and supply a subroutine `difeq`, called by `solvde`, that evaluates the  $\mathbf{s}$  matrix for each block.

Most of the arguments in the call to `solvde` have already been described, but some require discussion. Array  $y(j, k)$  contains the initial guess for the solution, with  $j$  labeling the dependent variables at mesh points  $k$ . The problem involves  $ne$  FDEs spanning points  $k=1, \dots, m$ .  $nb$  boundary conditions apply at the first point  $k=1$ . The array `indexv(j)` establishes the correspondence between columns of the  $\mathbf{s}$  matrix, equations (17.3.8), (17.3.10),

```

(a) D D D A A      V  A
     D D D A A      V  A
     D D D A A      V  A

(b) 1 0 0 S S      V  S
     0 1 0 S S      V  S
     0 0 1 S S      V  S
    
```

Figure 17.3.3. Reduction process for the first (upper left) block of the matrix in Figure 17.3.1. (a) Original form of the block, (b) final form. (See text for explanation.)

```

(a) 1 0 0 S S      V  S
     0 1 0 S S      V  S
     0 0 1 S S      V  S
     Z Z Z D D D D A A      V  A
     Z Z Z D D D D D A A      V  A
     Z Z Z D D D D D A A      V  A
     Z Z Z D D D D D A A      V  A
     Z Z Z D D D D D A A      V  A

(b) 1 0 0 S S      V  S
     0 1 0 S S      V  S
     0 0 1 S S      V  S
     0 0 0 1 0 0 0 0 S S      V  S
     0 0 0 0 1 0 0 0 S S      V  S
     0 0 0 0 0 1 0 0 S S      V  S
     0 0 0 0 0 0 1 0 S S      V  S
     0 0 0 0 0 0 0 1 S S      V  S
    
```

Figure 17.3.4. Reduction process for intermediate blocks of the matrix in Figure 17.3.1. (a) Original form, (b) final form. (See text for explanation.)

```

(a) 0 0 0 1 0 0 0 0 S S      V  S
     0 0 0 0 1 0 0 0 S S      V  S
     0 0 0 0 0 1 0 0 S S      V  S
     0 0 0 0 0 0 1 0 S S      V  S
     0 0 0 0 0 0 0 1 S S      V  S
           Z Z Z D D      V  A
           Z Z Z D D      V  A

(b) 0 0 0 1 0 0 0 0 S S      V  S
     0 0 0 0 1 0 0 0 S S      V  S
     0 0 0 0 0 1 0 0 S S      V  S
     0 0 0 0 0 0 1 0 S S      V  S
     0 0 0 0 0 0 0 1 S S      V  S
           0 0 0 1 0      V  S
           0 0 0 0 1      V  S
    
```

Figure 17.3.5. Reduction process for the last (lower right) block of the matrix in Figure 17.3.1. (a) Original form, (b) final form. (See text for explanation.)

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and (17.3.12), and the dependent variables. As described above it is essential that the  $nb$  boundary conditions at  $k=1$  involve the dependent variables referenced by the first  $nb$  columns of the  $s$  matrix. Thus, columns  $j$  of the  $s$  matrix can be ordered by the user in `difeq` to refer to derivatives with respect to the dependent variable `indexv(j)`.

The subroutine only attempts `itmax` correction cycles before returning, even if the solution has not converged. The parameters `conv`, `slowc`, `scalv` relate to convergence. Each inversion of the matrix produces corrections for `ne` variables at  $m$  mesh points. We want these to become vanishingly small as the iterations proceed, but we must define a measure for the size of corrections. This error “norm” is very problem specific, so the user might wish to rewrite this section of the code as appropriate. In the program below we compute a value for the average correction `err` by summing the absolute value of all corrections, weighted by a scale factor appropriate to each type of variable:

$$\text{err} = \frac{1}{m \times \text{ne}} \sum_{k=1}^m \sum_{j=1}^{\text{ne}} \frac{|\Delta Y(j,k)|}{\text{scalv}(j)} \quad (17.3.13)$$

When `err`  $\leq$  `conv`, the method has converged. Note that the user gets to supply an array `scalv` which measures the typical size of each variable.

Obviously, if `err` is large, we are far from a solution, and perhaps it is a bad idea to believe that the corrections generated from a first-order Taylor series are accurate. The number `slowc` modulates application of corrections. After each iteration we apply only a fraction of the corrections found by matrix inversion:

$$Y(j,k) \rightarrow Y(j,k) + \frac{\text{slowc}}{\max(\text{slowc}, \text{err})} \Delta Y(j,k) \quad (17.3.14)$$

Thus, when `err`  $>$  `slowc` only a fraction of the corrections are used, but when `err`  $\leq$  `slowc` the entire correction gets applied.

The call statement also supplies `solvde` with the array `y(1:nyj,1:nyk)` containing the initial trial solution, and workspace arrays `c(1:nci,1:ncj,1:nck)`, `s(1:nsi,1:nsj)`. The array `c` is the blockbuster: It stores the unreduced elements of the matrix built up for the backsubstitution step. If there are  $m$  mesh points, then there will be `nck`  $= m+1$  blocks, each requiring `nci`  $= ne$  rows and `ncj`  $= ne - nb + 1$  columns. Although large, this is small compared with  $(ne \times m)^2$  elements required for the whole matrix if we did not break it into blocks.

We now describe the workings of the user-supplied subroutine `difeq`. The parameters of the subroutine are given by

```
SUBROUTINE difeq(k,k1,k2,jsf,is1,isf,indexv,ne,s,nsi,nsj,y,nyj,nyk)
```

The only information returned from `difeq` to `solvde` is the matrix of derivatives `s(i,j)`; all other arguments are input to `difeq` and should not be altered. `k` indicates the current mesh point, or block number. `k1`, `k2` label the first and last point in the mesh. If `k`  $= k1$  or `k`  $>$  `k2`, the block involves the boundary conditions at the first or final points; otherwise the block acts on FDEs coupling variables at points `k-1`, `k`.

The convention on storing information into the array `s(i,j)` follows that used in equations (17.3.8), (17.3.10), and (17.3.12): Rows `i` label equations, columns `j` refer to derivatives with respect to dependent variables in the solution. Recall that each equation will depend on the `ne` dependent variables at either one or two points. Thus, `j` runs from 1 to either `ne` or  $2 \times ne$ . The column ordering for dependent variables at each point must agree with the list supplied in `indexv(j)`. Thus, for a block not at a boundary, the first column multiplies  $\Delta Y(1 = \text{indexv}(1), k-1)$ , and the column `ne+1` multiplies  $\Delta Y(1 = \text{indexv}(1), k)$ . `is1`, `isf` give the numbers of the starting and final rows that need to be filled in the  $s$  matrix for this block. `jsf` labels the column in which the difference equations  $E_{j,k}$  of equations (17.3.3)–(17.3.5) are stored. Thus,  $-s(i, jsf)$  is the vector on the right-hand side of the matrix. The reason for the minus sign is that `difeq` supplies the actual difference equation,  $E_{j,k}$ , not its negative. Note that `solvde` supplies a value for `jsf` such that the difference equation is put in the column *just after* all derivatives in the  $s$  matrix. Thus, `difeq` expects to find values entered into `s(i,j)` for rows `is1`  $\leq i \leq isf$  and  $1 \leq j \leq jsf$ .

Finally,  $s(1:nsi, 1:nsj)$  and  $y(1:nyj, 1:nyk)$  supply  $difeq$  with storage for  $s$  and the solution variables  $y$  for this iteration. An example of how to use this routine is given in the next section.

Many ideas in the following code are due to Eggleton [1].

```

SUBROUTINE solvde(itmax,conv,slowc,scalv,indexv,ne,nb,m,
*   y,nyj,nyk,c,nci,ncj,nck,s,nsi,nsj)
INTEGER itmax,m,nb,nci,ncj,nck,ne,nsi,nsj,
*   nyj,nyk,indexv(nyj),NMAX
REAL conv,slowc,c(nci,ncj,nck),s(nsi,nsj),
*   scalv(nyj),y(nyj,nyk)
PARAMETER (NMAX=10)           Largest expected value of ne.
C  USES bksub,difeq,pinvs,red
   Driver routine for solution of two point boundary value problems by relaxation. itmax is the
   maximum number of iterations. conv is the convergence criterion (see text). slowc con-
   trols the fraction of corrections actually used after each iteration. scalv(1:nyj) con-
   tains typical sizes for each dependent variable, used to weight errors. indexv(1:nyj) lists the
   column ordering of variables used to construct the matrix s of derivatives. (The nb boundary
   conditions at the first mesh point must contain some dependence on the first nb variables
   listed in indexv.) The problem involves ne equations for ne adjustable dependent variables
   at each point. At the first mesh point there are nb boundary conditions. There are a total
   of m mesh points. y(1:nyj, 1:nyk) is the two-dimensional array that contains the initial
   guess for all the dependent variables at each mesh point. On each iteration, it is updated by
   the calculated correction. The arrays c(1:nci, 1:ncj, 1:nck), s(1:nsi, 1:nsj) sup-
   ply dummy storage used by the relaxation code; the minimum dimensions must satisfy:
   nci=ne, ncj=ne-nb+1, nck=m+1, nsi=ne, nsj=2*ne+1.
INTEGER ic1,ic2,ic3,ic4,it,j,j1,j2,j3,j4,j5,j6,j7,j8,
*   j9,jc1,jcf,jv,k,k1,k2,km,kp,nvars,kmax(NMAX)
REAL err,errj,fac,vmax,vz,ermax(NMAX)
k1=1           Set up row and column markers.
k2=m
nvars=ne*m
j1=1
j2=nb
j3=nb+1
j4=ne
j5=j4+j1
j6=j4+j2
j7=j4+j3
j8=j4+j4
j9=j8+j1
ic1=1
ic2=ne-nb
ic3=ic2+1
ic4=ne
jc1=1
jcf=ic3
do 16 it=1,itmax           Primary iteration loop.
  k=k1           Boundary conditions at first point.
  call difeq(k,k1,k2,j9,ic3,ic4,indexv,ne,s,nsi,nsj,y,nyj,nyk)
  call pinvs(ic3,ic4,j5,j9,jc1,k1,c,nci,ncj,nck,s,nsi,nsj)
  do 11 k=k1+1,k2           Finite difference equations at all point pairs.
    kp=k-1
    call difeq(k,k1,k2,j9,ic1,ic4,indexv,ne,s,nsi,nsj,y,nyj,nyk)
    call red(ic1,ic4,j1,j2,j3,j4,j9,ic3,jc1,jcf,kp,
*     c,nci,ncj,nck,s,nsi,nsj)
    call pinvs(ic1,ic4,j3,j9,jc1,k,c,nci,ncj,nck,s,nsi,nsj)
  enddo 11
  k=k2+1           Final boundary conditions.
  call difeq(k,k1,k2,j9,ic1,ic2,indexv,ne,s,nsi,nsj,y,nyj,nyk)
  call red(ic1,ic2,j5,j6,j7,j8,j9,ic3,jc1,jcf,k2,
*   c,nci,ncj,nck,s,nsi,nsj)
  call pinvs(ic1,ic2,j7,j9,jcf,k2+1,c,nci,ncj,nck,s,nsi,nsj)

```

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```

call bksub(ne,nb,jcf,k1,k2,c,nci,ncj,nck)      Backsubstitution.
err=0.
do 13 j=1,ne          Convergence check, accumulate average error.
  jv=indexv(j)
  errj=0.
  km=0
  vmax=0.
  do 12 k=k1,k2      Find point with largest error, for each dependent variable.
    vz=abs(c(jv,1,k))
    if(vz.gt.vmax) then
      vmax=vz
      km=k
    endif
    errj=errj+vz
  enddo 12
  err=err+errj/scalv(j)      Note weighting for each dependent variable.
  ermax(j)=c(jv,1,km)/scalv(j)
  kmax(j)=km
enddo 13
err=err/nvars
fac=slowc/max(slowc,err)      Reduce correction applied when error is large.
do 15 j=1,ne          Apply corrections.
  jv=indexv(j)
  do 14 k=k1,k2
    y(j,k)=y(j,k)-fac*c(jv,1,k)
  enddo 14
enddo 15
write(*,100) it,err,fac      Summary of corrections for this step. Point with largest
if(err.lt.conv) return      error for each variable can be monitored by writ-
ing out kmax and ermax.
enddo 16
100 pause 'itmax exceeded in solvde'      Convergence failed.
format(1x,i4,2f12.6)
return
END
SUBROUTINE bksub(ne,nb,jf,k1,k2,c,nci,ncj,nck)
INTEGER jf,k1,k2,nb,nci,ncj,nck,ne
REAL c(nci,ncj,nck)
Backsubstitution, used internally by solvde.
INTEGER i,im,j,k,kp,nbf
REAL xx
nbf=ne-nb
im=1
do 13 k=k2,k1,-1      Use recurrence relations to eliminate remaining dependences.
  if (k.eq.k1) im=nbf+1      Special handling of first point.
  kp=k+1
  do 12 j=1,nbf
    xx=c(j,jf,kp)
    do 11 i=im,ne
      c(i,jf,k)=c(i,jf,k)-c(i,j,k)*xx
    enddo 11
  enddo 12
enddo 13
do 16 k=k1,k2          Reorder corrections to be in column 1.
  kp=k+1
  do 14 i=1,nb
    c(i,1,k)=c(i+nbf,jf,k)
  enddo 14
  do 15 i=1,nbf
    c(i+nbf,1,k)=c(i,jf,kp)
  enddo 15
enddo 16
return
END

```

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```

SUBROUTINE pinvs(ie1,ie2,je1,jsf,jc1,k,c,nci,ncj,nck,s,nsi,nsj)
INTEGER ie1,ie2,jc1,je1,jsf,k,nci,ncj,nck,nsi,nsj,NMAX
REAL c(nci,ncj,nck),s(nsi,nsj)
PARAMETER (NMAX=10)
    Diagonalize the square subsection of the s matrix, and store the recursion coefficients in
    c; used internally by solvde.
INTEGER i,icoff,id,ipiv,irow,j,jcoff,je2,jp,jpiv,js1,indxr(NMAX)
REAL big,dum,piv,pivinv,pscl(NMAX)
je2=je1+ie2-ie1
js1=je2+1
do 12 i=ie1,ie2          Implicit pivoting, as in §2.1.
    big=0.
    do 11 j=je1,je2
        if(abs(s(i,j)).gt.big) big=abs(s(i,j))
    enddo 11
    if(big.eq.0.) pause 'singular matrix, row all 0 in pinvs'
    pscl(i)=1./big
    indxr(i)=0
enddo 12
do 18 id=ie1,ie2
    piv=0.
    do 14 i=ie1,ie2      Find pivot element.
        if(indxr(i).eq.0) then
            big=0.
            do 13 j=je1,je2
                if(abs(s(i,j)).gt.big) then
                    jp=j
                    big=abs(s(i,j))
                endif
            enddo 13
            if(big*pscl(i).gt.piv) then
                ipiv=i
                jpiv=jp
                piv=big*pscl(i)
            endif
        endif
    enddo 14
    if(s(ipiv,jpiv).eq.0.) pause 'singular matrix in pinvs'
    indxr(ipiv)=jpiv      In place reduction. Save column ordering.
    pivinv=1./s(ipiv,jpiv)
    do 15 j=je1,jsf      Normalize pivot row.
        s(ipiv,j)=s(ipiv,j)*pivinv
    enddo 15
    s(ipiv,jpiv)=1.
    do 17 i=ie1,ie2      Reduce nonpivot elements in column.
        if(indxr(i).ne.jpiv) then
            if(s(i,jpiv).ne.0.) then
                dum=s(i,jpiv)
                do 16 j=je1,jsf
                    s(i,j)=s(i,j)-dum*s(ipiv,j)
                enddo 16
                s(i,jpiv)=0.
            endif
        endif
    enddo 17
enddo 18
jcoff=jc1-js1          Sort and store unreduced coefficients.
icoff=ie1-je1
do 21 i=ie1,ie2
    irow=indxr(i)+icoff
    do 19 j=js1,jsf
        c(irow,j+jcoff,k)=s(i,j)
    enddo 19
enddo 21

```

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```

return
END

SUBROUTINE red(iz1,iz2,jz1,jz2,jm1,jm2,jmf,ic1,jc1,jcf,kc,
*      c,nci,ncj,nck,s,nsi,nsj)
INTEGER ic1,iz1,iz2,jc1,jcf,jm1,jm2,jmf,jz1,jz2,kc,nci,ncj,
*      nck,nsi,nsj
REAL c(nci,ncj,nck),s(nsi,nsj)
  Reduce columns jz1-jz2 of the s matrix, using previous results as stored in the c matrix.
  Only columns jm1-jm2,jmf are affected by the prior results. red is used internally by
  solvde.
INTEGER i,ic,j,l,loff
REAL vx
loff=jc1-jm1
ic=ic1
do 14 j=jz1,jz2          Loop over columns to be zeroed.
  do 12 l=jm1,jm2        Loop over columns altered.
    vx=c(ic,l+loff,kc)
    do 11 i=iz1,iz2      Loop over rows.
      s(i,l)=s(i,l)-s(i,j)*vx
    enddo 11
  enddo 12
  vx=c(ic,jcf,kc)
  do 13 i=iz1,iz2        Plus final element.
    s(i,jmf)=s(i,jmf)-s(i,j)*vx
  enddo 13
  ic=ic+1
enddo 14
return
END

```

## “Algebraically Difficult” Sets of Differential Equations

Relaxation methods allow you to take advantage of an additional opportunity that, while not obvious, can speed up some calculations enormously. It is not necessary that the set of variables  $y_{j,k}$  correspond exactly with the dependent variables of the original differential equations. They can be related to those variables through algebraic equations. Obviously, it is necessary only that the solution variables allow us to *evaluate* the functions  $y, g, \mathbf{B}, \mathbf{C}$  that are used to construct the FDEs from the ODEs. In some problems  $g$  depends on functions of  $y$  that are known only implicitly, so that iterative solutions are necessary to evaluate functions in the ODEs. Often one can dispense with this “internal” nonlinear problem by defining a new set of variables from which both  $y, g$  and the boundary conditions can be obtained directly. A typical example occurs in physical problems where the equations require solution of a complex equation of state that can be expressed in more convenient terms using variables other than the original dependent variables in the ODE. While this approach is analogous to performing an *analytic* change of variables directly on the original ODEs, such an analytic transformation might be prohibitively complicated. The change of variables in the relaxation method is easy and requires no analytic manipulations.

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- Eggleton, P.P. 1971, *Monthly Notices of the Royal Astronomical Society*, vol. 151, pp. 351–364. [1]  
 Keller, H.B. 1968, *Numerical Methods for Two-Point Boundary-Value Problems* (Waltham, MA: Blaisdell).  
 Kippenhan, R., Weigert, A., and Hofmeister, E. 1968, in *Methods in Computational Physics*, vol. 7 (New York: Academic Press), pp. 129ff.

## 17.4 A Worked Example: Spheroidal Harmonics

The best way to understand the algorithms of the previous sections is to see them employed to solve an actual problem. As a sample problem, we have selected the computation of spheroidal harmonics. (The more common name is spheroidal angle functions, but we prefer the explicit reminder of the kinship with spherical harmonics.) We will show how to find spheroidal harmonics, first by the method of relaxation (§17.3), and then by the methods of shooting (§17.1) and shooting to a fitting point (§17.2).

Spheroidal harmonics typically arise when certain partial differential equations are solved by separation of variables in spheroidal coordinates. They satisfy the following differential equation on the interval  $-1 \leq x \leq 1$ :

$$\frac{d}{dx} \left[ (1-x^2) \frac{dS}{dx} \right] + \left( \lambda - c^2 x^2 - \frac{m^2}{1-x^2} \right) S = 0 \quad (17.4.1)$$

Here  $m$  is an integer,  $c$  is the “oblateness parameter,” and  $\lambda$  is the eigenvalue. Despite the notation,  $c^2$  can be positive or negative. For  $c^2 > 0$  the functions are called “prolate,” while if  $c^2 < 0$  they are called “oblate.” The equation has singular points at  $x = \pm 1$  and is to be solved subject to the boundary conditions that the solution be regular at  $x = \pm 1$ . Only for certain values of  $\lambda$ , the eigenvalues, will this be possible.

If we consider first the spherical case, where  $c = 0$ , we recognize the differential equation for Legendre functions  $P_n^m(x)$ . In this case the eigenvalues are  $\lambda_{mn} = n(n+1)$ ,  $n = m, m+1, \dots$ . The integer  $n$  labels successive eigenvalues for fixed  $m$ : When  $n = m$  we have the lowest eigenvalue, and the corresponding eigenfunction has no nodes in the interval  $-1 < x < 1$ ; when  $n = m+1$  we have the next eigenvalue, and the eigenfunction has one node inside  $(-1, 1)$ ; and so on.

A similar situation holds for the general case  $c^2 \neq 0$ . We write the eigenvalues of (17.4.1) as  $\lambda_{mn}(c)$  and the eigenfunctions as  $S_{mn}(x; c)$ . For fixed  $m$ ,  $n = m, m+1, \dots$  labels the successive eigenvalues.

The computation of  $\lambda_{mn}(c)$  and  $S_{mn}(x; c)$  traditionally has been quite difficult. Complicated recurrence relations, power series expansions, etc., can be found in references [1-3]. Cheap computing makes evaluation by direct solution of the differential equation quite feasible.

The first step is to investigate the behavior of the solution near the singular points  $x = \pm 1$ . Substituting a power series expansion of the form

$$S = (1 \pm x)^\alpha \sum_{k=0}^{\infty} a_k (1 \pm x)^k \quad (17.4.2)$$

in equation (17.4.1), we find that the regular solution has  $\alpha = m/2$ . (Without loss of generality we can take  $m \geq 0$  since  $m \rightarrow -m$  is a symmetry of the equation.) We get an equation that is numerically more tractable if we factor out this behavior. Accordingly we set

$$S = (1-x^2)^{m/2} y \quad (17.4.3)$$

We then find from (17.4.1) that  $y$  satisfies the equation

$$(1-x^2) \frac{d^2 y}{dx^2} - 2(m+1)x \frac{dy}{dx} + (\mu - c^2 x^2)y = 0 \quad (17.4.4)$$

where

$$\mu \equiv \lambda - m(m + 1) \tag{17.4.5}$$

Both equations (17.4.1) and (17.4.4) are invariant under the replacement  $x \rightarrow -x$ . Thus the functions  $S$  and  $y$  must also be invariant, except possibly for an overall scale factor. (Since the equations are linear, a constant multiple of a solution is also a solution.) Because the solutions will be normalized, the scale factor can only be  $\pm 1$ . If  $n - m$  is odd, there are an odd number of zeros in the interval  $(-1, 1)$ . Thus we must choose the antisymmetric solution  $y(-x) = -y(x)$  which has a zero at  $x = 0$ . Conversely, if  $n - m$  is even we must have the symmetric solution. Thus

$$y_{mn}(-x) = (-1)^{n-m} y_{mn}(x) \tag{17.4.6}$$

and similarly for  $S_{mn}$ .

The boundary conditions on (17.4.4) require that  $y$  be regular at  $x = \pm 1$ . In other words, near the endpoints the solution takes the form

$$y = a_0 + a_1(1 - x^2) + a_2(1 - x^2)^2 + \dots \tag{17.4.7}$$

Substituting this expansion in equation (17.4.4) and letting  $x \rightarrow 1$ , we find that

$$a_1 = -\frac{\mu - c^2}{4(m + 1)} a_0 \tag{17.4.8}$$

Equivalently,

$$y'(1) = \frac{\mu - c^2}{2(m + 1)} y(1) \tag{17.4.9}$$

A similar equation holds at  $x = -1$  with a minus sign on the right-hand side. The irregular solution has a different relation between function and derivative at the endpoints.

Instead of integrating the equation from  $-1$  to  $1$ , we can exploit the symmetry (17.4.6) to integrate from  $0$  to  $1$ . The boundary condition at  $x = 0$  is

$$\begin{aligned} y(0) &= 0, & n - m \text{ odd} \\ y'(0) &= 0, & n - m \text{ even} \end{aligned} \tag{17.4.10}$$

A third boundary condition comes from the fact that any constant multiple of a solution  $y$  is a solution. We can thus *normalize* the solution. We adopt the normalization that the function  $S_{mn}$  has the same limiting behavior as  $P_n^m$  at  $x = 1$ :

$$\lim_{x \rightarrow 1} (1 - x^2)^{-m/2} S_{mn}(x; c) = \lim_{x \rightarrow 1} (1 - x^2)^{-m/2} P_n^m(x) \tag{17.4.11}$$

Various normalization conventions in the literature are tabulated by Flammer [1].

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Imposing three boundary conditions for the second-order equation (17.4.4) turns it into an eigenvalue problem for  $\lambda$  or equivalently for  $\mu$ . We write it in the standard form by setting

$$y_1 = y \quad (17.4.12)$$

$$y_2 = y' \quad (17.4.13)$$

$$y_3 = \mu \quad (17.4.14)$$

Then

$$y_1' = y_2 \quad (17.4.15)$$

$$y_2' = \frac{1}{1-x^2} [2x(m+1)y_2 - (y_3 - c^2x^2)y_1] \quad (17.4.16)$$

$$y_3' = 0 \quad (17.4.17)$$

The boundary condition at  $x = 0$  in this notation is

$$\begin{aligned} y_1 &= 0, & n - m & \text{ odd} \\ y_2 &= 0, & n - m & \text{ even} \end{aligned} \quad (17.4.18)$$

At  $x = 1$  we have two conditions:

$$y_2 = \frac{y_3 - c^2}{2(m+1)} y_1 \quad (17.4.19)$$

$$y_1 = \lim_{x \rightarrow 1} (1-x^2)^{-m/2} P_n^m(x) = \frac{(-1)^m (n+m)!}{2^m m! (n-m)!} \equiv \gamma \quad (17.4.20)$$

We are now ready to illustrate the use of the methods of previous sections on this problem.

### Relaxation

If we just want a few isolated values of  $\lambda$  or  $S$ , shooting is probably the quickest method. However, if we want values for a large sequence of values of  $c$ , relaxation is better. Relaxation rewards a good initial guess with rapid convergence, and the previous solution should be a good initial guess if  $c$  is changed only slightly.

For simplicity, we choose a uniform grid on the interval  $0 \leq x \leq 1$ . For a total of  $M$  mesh points, we have

$$h = \frac{1}{M-1} \quad (17.4.21)$$

$$x_k = (k-1)h, \quad k = 1, 2, \dots, M \quad (17.4.22)$$

At interior points  $k = 2, 3, \dots, M$ , equation (17.4.15) gives

$$E_{1,k} = y_{1,k} - y_{1,k-1} - \frac{h}{2}(y_{2,k} + y_{2,k-1}) \quad (17.4.23)$$

Equation (17.4.16) gives

$$E_{2,k} = y_{2,k} - y_{2,k-1} - \beta_k \times \left[ \frac{(x_k + x_{k-1})(m+1)(y_{2,k} + y_{2,k-1})}{2} - \alpha_k \frac{(y_{1,k} + y_{1,k-1})}{2} \right] \quad (17.4.24)$$

where

$$\alpha_k = \frac{y_{3,k} + y_{3,k-1}}{2} - \frac{c^2(x_k + x_{k-1})^2}{4} \quad (17.4.25)$$

$$\beta_k = \frac{h}{1 - \frac{1}{4}(x_k + x_{k-1})^2} \quad (17.4.26)$$

Finally, equation (17.4.17) gives

$$E_{3,k} = y_{3,k} - y_{3,k-1} \quad (17.4.27)$$

Now recall that the matrix of partial derivatives  $S_{i,j}$  of equation (17.3.8) is defined so that  $i$  labels the equation and  $j$  the variable. In our case,  $j$  runs from 1 to 3 for  $y_j$  at  $k - 1$  and from 4 to 6 for  $y_j$  at  $k$ . Thus equation (17.4.23) gives

$$\begin{aligned} S_{1,1} &= -1, & S_{1,2} &= -\frac{h}{2}, & S_{1,3} &= 0 \\ S_{1,4} &= 1, & S_{1,5} &= -\frac{h}{2}, & S_{1,6} &= 0 \end{aligned} \quad (17.4.28)$$

Similarly equation (17.4.24) yields

$$\begin{aligned} S_{2,1} &= \alpha_k \beta_k / 2, & S_{2,2} &= -1 - \beta_k(x_k + x_{k-1})(m+1)/2, \\ S_{2,3} &= \beta_k(y_{1,k} + y_{1,k-1})/4 & S_{2,4} &= S_{2,1}, \\ S_{2,5} &= 2 + S_{2,2}, & S_{2,6} &= S_{2,3} \end{aligned} \quad (17.4.29)$$

while from equation (17.4.27) we find

$$\begin{aligned} S_{3,1} &= 0, & S_{3,2} &= 0, & S_{3,3} &= -1 \\ S_{3,4} &= 0, & S_{3,5} &= 0, & S_{3,6} &= 1 \end{aligned} \quad (17.4.30)$$

At  $x = 0$  we have the boundary condition

$$E_{3,1} = \begin{cases} y_{1,1}, & n - m \text{ odd} \\ y_{2,1}, & n - m \text{ even} \end{cases} \quad (17.4.31)$$

Recall the convention adopted in the `solvde` routine that for one boundary condition at  $k = 1$  only  $S_{3,j}$  can be nonzero. Also,  $j$  takes on the values 4 to 6 since the boundary condition involves only  $y_k$ , not  $y_{k-1}$ . Accordingly, the only nonzero values of  $S_{3,j}$  at  $x = 0$  are

$$\begin{aligned} S_{3,4} &= 1, & n - m \text{ odd} \\ S_{3,5} &= 1, & n - m \text{ even} \end{aligned} \quad (17.4.32)$$

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At  $x = 1$  we have

$$E_{1,M+1} = y_{2,M} - \frac{y_{3,M} - c^2}{2(m+1)} y_{1,M} \quad (17.4.33)$$

$$E_{2,M+1} = y_{1,M} - \gamma \quad (17.4.34)$$

Thus

$$S_{1,4} = -\frac{y_{3,M} - c^2}{2(m+1)}, \quad S_{1,5} = 1, \quad S_{1,6} = -\frac{y_{1,M}}{2(m+1)} \quad (17.4.35)$$

$$S_{2,4} = 1, \quad S_{2,5} = 0, \quad S_{2,6} = 0 \quad (17.4.36)$$

Here now is the sample program that implements the above algorithm. We need a main program, `sfroid`, that calls the routine `solvde`, and we must supply the subroutine `difeq` called by `solvde`. For simplicity we choose an equally spaced mesh of  $m = 41$  points, that is,  $h = .025$ . As we shall see, this gives good accuracy for the eigenvalues up to moderate values of  $n - m$ .

Since the boundary condition at  $x = 0$  does not involve  $y_1$  if  $n - m$  is even, we have to use the `indexv` feature of `solvde`. Recall that the value of `indexv(j)` describes which column of  $s(i, j)$  the variable  $y(j)$  has been put in. If  $n - m$  is even, we need to interchange the columns for  $y_1$  and  $y_2$  so that there is not a zero pivot element in  $s(i, j)$ .

The program prompts for values of  $m$  and  $n$ . It then computes an initial guess for  $y$  based on the Legendre function  $P_n^m$ . It next prompts for  $c^2$ , solves for  $y$ , prompts for  $c^2$ , solves for  $y$  using the previous values as an initial guess, and so on.

```

PROGRAM sfroid
INTEGER NE,M,NB,NCI,NCJ,NCK,NSI,NSJ,NYJ,NYK
COMMON /sfrcom/ x,h,mm,n,c2,anorm          Communicates with difeq.
PARAMETER (NE=3,M=41,NB=1,NCI=NE,NCJ=NE-NB+1,NCK=M+1,NSI=NE,
*       NSJ=2*NE+1,NYJ=NE,NYK=M)
C  USES plgndr,solvde
    Sample program using solvde. Computes eigenvalues of spheroidal harmonics  $S_{mn}(x; c)$ 
    for  $m \geq 0$  and  $n \geq m$ . In the program,  $m$  is mm,  $c^2$  is c2, and  $\gamma$  of equation (17.4.20)
    is anorm.
INTEGER i,itmax,k,mm,n,indexv(NE)
REAL anorm,c2,conv,deriv,fac1,fac2,h,q1,slowc,
*   c(NCI,NCJ,NCK),s(NSI,NSJ),scalv(NE),x(M),y(NE,M),plgndr
itmax=100
conv=5.e-6
slowc=1.
h=1./(M-1)
c2=0.
write(*,*)'ENTER M,N'
read(*,*)mm,n
if(mod(n+mm,2).eq.1)then                No interchanges necessary.
    indexv(1)=1
    indexv(2)=2
    indexv(3)=3
else                                      Interchange  $y_1$  and  $y_2$ .
    indexv(1)=2
    indexv(2)=1
    indexv(3)=3
endif
anorm=1.                                Compute  $\gamma$ .

```



```

if (mm.NE.0) then
  q1=n
  do 11 i=1,mm
    anorm=-.5*anorm*(n+i)*(q1/i)
    q1=q1-1.
  enddo 11
endif
do 12 k=1,M-1
  x(k)=(k-1)*h
  fac1=1.-x(k)**2
  fac2=fac1**(-mm/2.)
  y(1,k)=plgndr(n,mm,x(k))*fac2
  deriv=-((n-mm+1)*plgndr(n+1,mm,x(k))-(n+1)*
  * x(k)*plgndr(n,mm,x(k)))/fac1
  y(2,k)=mm*x(k)*y(1,k)/fac1+deriv*fac2
  y(3,k)=n*(n+1)-mm*(mm+1)
enddo 12
x(M)=1.
y(1,M)=anorm
y(3,M)=n*(n+1)-mm*(mm+1)
y(2,M)=(y(3,M)-c2)*y(1,M)/(2.*(mm+1.))
scalv(1)=abs(anorm)
scalv(2)=max(abs(anorm),y(2,M))
scalv(3)=max(1.,y(3,M))
1 continue
write (*,*) 'ENTER C**2 OR 999 TO END'
read (*,*) c2
if (c2.eq.999.) stop
call solvde(itmax,conv,slowc,scalv,indexv,NE,NB,M,y,NYJ,NYK,
* c,NCI,NCJ,NCK,s,NSI,NSJ)
write (*,*) ' M = ',mm,' N = ',n,
* ' C**2 = ',c2,' LAMBDA = ',y(3,1)+mm*(mm+1)
goto 1
END

```

Initial guess.

$P_n^m$  from §6.8.

Derivative of  $P_n^m$  from a recurrence relation.

Initial guess at  $x = 1$  done separately.

for another value of  $c^2$ .

```

SUBROUTINE difeq(k,k1,k2,jsf,is1,isf,indexv,ne,s,nsi,nsj,y,nyj,nyk)
INTEGER is1,isf,jsf,k,k1,k2,ne,nsi,nsj,nyj,nyk,indexv(nyj),M
REAL s(nsi,nsj),y(nyj,nyk)
COMMON /sfrcom/ x,h,mm,n,c2,anorm
PARAMETER (M=41)
  Returns matrix s(i,j) for solvde.
INTEGER mm,n
REAL anorm,c2,h,temp,temp2,x(M)
if(k.eq.k1) then
  if(mod(n+mm,2).eq.1) then
    s(3,3+indexv(1))=1.
    s(3,3+indexv(2))=0.
    s(3,3+indexv(3))=0.
    s(3,jsf)=y(1,1)
  else
    s(3,3+indexv(1))=0.
    s(3,3+indexv(2))=1.
    s(3,3+indexv(3))=0.
    s(3,jsf)=y(2,1)
  endif
else if(k.gt.k2) then
  s(1,3+indexv(1))=-(y(3,M)-c2)/(2.*(mm+1.))
  s(1,3+indexv(2))=1.
  s(1,3+indexv(3))=-y(1,M)/(2.*(mm+1.))
  s(1,jsf)=y(2,M)-(y(3,M)-c2)*y(1,M)/(2.*(mm+1.))
  s(2,3+indexv(1))=1.
  s(2,3+indexv(2))=0.

```

Boundary condition at first point.

Equation (17.4.32).

Equation (17.4.31).

Equation (17.4.32).

Equation (17.4.31).

Boundary conditions at last point.

Equation (17.4.35).

Equation (17.4.33).

Equation (17.4.36).

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```

s(2,3+indexv(3))=0.
s(2,jsf)=y(1,M)-anorm      Equation (17.4.34).
else                          Interior point.
s(1,indexv(1))=-1.          Equation (17.4.28).
s(1,indexv(2))=-.5*h
s(1,indexv(3))=0.
s(1,3+indexv(1))=1.
s(1,3+indexv(2))=-.5*h
s(1,3+indexv(3))=0.
temp=h/(1.-(x(k)+x(k-1))**2*.25)
temp2=.5*(y(3,k)+y(3,k-1))-c2*.25*(x(k)+x(k-1))**2
s(2,indexv(1))=temp*temp2*.5      Equation (17.4.29).
s(2,indexv(2))=-1.-.5*temp*(mm+1.)*(x(k)+x(k-1))
s(2,indexv(3))=.25*temp*(y(1,k)+y(1,k-1))
s(2,3+indexv(1))=s(2,indexv(1))
s(2,3+indexv(2))=2.+s(2,indexv(2))
s(2,3+indexv(3))=s(2,indexv(3))
s(3,indexv(1))=0.              Equation (17.4.30).
s(3,indexv(2))=0.
s(3,indexv(3))=-1.
s(3,3+indexv(1))=0.
s(3,3+indexv(2))=0.
s(3,3+indexv(3))=1.
s(1,jsf)=y(1,k)-y(1,k-1)-.5*h*(y(2,k)+y(2,k-1))      Equation (17.4.23).
s(2,jsf)=y(2,k)-y(2,k-1)-temp*(x(k)+x(k-1))          Equation (17.4.24).
*      *.5*(mm+1.)*(y(2,k)+y(2,k-1))-temp2*
*      *.5*(y(1,k)+y(1,k-1))
s(3,jsf)=y(3,k)-y(3,k-1)      Equation (17.4.27).
endif
return
END

```

You can run the program and check it against values of  $\lambda_{mn}(c)$  given in the tables at the back of Flammer's book [1] or in Table 21.1 of Abramowitz and Stegun [2]. Typically it converges in about 3 iterations. The table below gives a few comparisons.

Selected Output of sfroid				
$m$	$n$	$c^2$	$\lambda_{\text{exact}}$	$\lambda_{\text{sfroid}}$
2	2	0.1	6.01427	6.01427
		1.0	6.14095	6.14095
		4.0	6.54250	6.54253
2	5	1.0	30.4361	30.4372
		16.0	36.9963	37.0135
4	11	-1.0	131.560	131.554

## Shooting

To solve the same problem via shooting (§17.1), we supply a subroutine `derivs` that implements equations (17.4.15)–(17.4.17). We will integrate the equations over the range  $-1 \leq x \leq 0$ . We provide the subroutine `load` which sets the eigenvalue  $y_3$  to its current best estimate,  $v(1)$ . It also sets the boundary values of  $y_1$  and

$y_2$  using equations (17.4.20) and (17.4.19) (with a minus sign corresponding to  $x = -1$ ). Note that the boundary condition is actually applied a distance  $dx$  from the boundary to avoid having to evaluate  $y'_2$  right on the boundary. The subroutine score follows from equation (17.4.18).

```

PROGRAM sphoot
  Sample program using shoot. Computes eigenvalues of spheroidal harmonics  $S_{mn}(x; c)$  for
   $m \geq 0$  and  $n \geq m$ . Be sure that routine funcv for newt is provided by shoot (§17.1).
  INTEGER i,m,n,nvar,N2
  PARAMETER (N2=1)
  REAL c2,dx,gamma,q1,x1,x2,v(N2)
  LOGICAL check
  COMMON /sphcom/ c2,gamma,dx,m,n      Communicates with load, score, and derivs.
  COMMON /caller/ x1,x2,nvar          Communicates with shoot.
C  USES newt
  dx=1.e-4                          Avoid evaluating derivatives exactly at  $x = -1$ .
  nvar=3                             Number of equations.
1  write(*,*) 'input m,n,c-squared (999 to end)'
  read(*,*) m,n,c2
  if (c2.eq.999.) stop
  if ((n.lt.m).or.(m.lt.0)) goto 1
  gamma=1.0                          Compute  $\gamma$  of equation (17.4.20).
  q1=n
  do 11 i=1,m
    gamma=-0.5*gamma*(n+i)*(q1/i)
    q1=q1-1.0
  enddo 11
  v(1)=n*(n+1)-m*(m+1)+c2/2.0        Initial guess for eigenvalue.
  x1=-1.0+dx                          Set range of integration.
  x2=0.0
  call newt(v,N2,check)                Find v that zeros function f in score.
  if(check)then
    write(*,*) 'shoot failed; bad initial guess'
  else
    write(*,'(1x,t6,a)') 'mu(m,n)'
    write(*,'(1x,f12.6)') v(1)
    goto 1
  endif
END

SUBROUTINE load(x1,v,y)
  INTEGER m,n
  REAL c2,dx,gamma,x1,y1,v(1),y(3)
  COMMON /sphcom/ c2,gamma,dx,m,n
  Supplies starting values for integration at  $x = -1 + dx$ .
  y(3)=v(1)
  if(mod(n-m,2).eq.0)then
    y1=gamma
  else
    y1=-gamma
  endif
  y(2)=-y(3)-c2*y1/(2*(m+1))
  y(1)=y1+y(2)*dx
  return
END

SUBROUTINE score(x2,y,f)
  INTEGER m,n
  REAL c2,dx,gamma,x2,f(1),y(3)
  COMMON /sphcom/ c2,gamma,dx,m,n
  Tests whether boundary condition at  $x = 0$  is satisfied.
  if (mod(n-m,2).eq.0) then
    f(1)=y(2)

```

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```

else
  f(1)=y(1)
endif
return
END

SUBROUTINE derivs(x,y,dydx)
INTEGER m,n
REAL c2,dx,gamma,x,dydx(3),y(3)
COMMON /sphcom/ c2,gamma,dx,m,n
  Evaluates derivatives for odeint.
dydx(1)=y(2)
dydx(2)=(2.0*x*(m+1.0)*y(2)-(y(3)-c2*x*x)*y(1))/(1.0-x*x)
dydx(3)=0.0
return
END

```

### Shooting to a Fitting Point

For variety we illustrate `shootf` from §17.2 by integrating over the whole range  $-1 + dx \leq x \leq 1 - dx$ , with the fitting point chosen to be at  $x = 0$ . The routine `derivs` is identical to the one for `shoot`. Now, however, there are two load routines. The routine `load1` for  $x = -1$  is essentially identical to `load` above. At  $x = 1$ , `load2` sets the function value  $y_1$  and the eigenvalue  $y_3$  to their best current estimates, `v2(1)` and `v2(2)`, respectively. If you quite sensibly make your initial guess of the eigenvalue the same in the two intervals, then `v1(1)` will stay equal to `v2(2)` during the iteration. The subroutine `score` simply checks whether all three function values match at the fitting point.

```

PROGRAM sphfpt
  Sample program using shootf. Computes eigenvalues of spheroidal harmonics  $S_{mn}(x; c)$ 
  for  $m \geq 0$  and  $n \geq m$ . Be sure that routine funcv for newt is provided by shootf (§17.2).
  The routine derivs is the same as for sphoot.
INTEGER i,m,n,nvar,nn2,N1,N2,NTOT
REAL DXX
PARAMETER (N1=2,N2=1,NTOT=N1+N2,DXX=1.e-4)
REAL c2,dx,gamma,q1,x1,x2,xf,v1(N2),v2(N1),v(NTOT)
LOGICAL check
COMMON /sphcom/ c2,gamma,dx,m,n
  Communicates with load1, load2, score, and derivs.
COMMON /caller/ x1,x2,xf,nvar,nn2      Communicates with shootf.
EQUIVALENCE (v1(1),v(1)),(v2(1),v(N2+1))
C  USES newt
nvar=NTOT          Number of equations.
nn2=N2
dx=DXX            Avoid evaluating derivatives exactly at  $x = \pm 1$ .
1 write(*,*) 'input m,n,c-squared (999 to end)'
  read(*,*) m,n,c2
  if (c2.eq.999.) stop
  if ((n.lt.m).or.(m.lt.0)) goto 1
  gamma=1.0       Compute  $\gamma$  of equation (17.4.20).
  q1=n
  do 11 i=1,m
    gamma=-0.5*gamma*(n+i)*(q1/i)
    q1=q1-1.0
  enddo 11
  v1(1)=n*(n+1)-m*(m+1)+c2/2.0      Initial guess for eigenvalue and function value.
  v2(2)=v1(1)

```

```

v2(1)=gamma*(1.-(v2(2)-c2)*dx/(2*(m+1)))
x1=-1.0+dx           Set range of integration.
x2=1.0-dx
xf=0.                Fitting point.
call newt(v,NTOT,check) Find v that zeros function f in score.
if(check)then
  write(*,*)'shootf failed; bad initial guess'
else
  write(*,'(1x,t6,a)') 'mu(m,n)'
  write(*,'(1x,f12.6)') v1(1)
  goto 1
endif
END

SUBROUTINE load1(x1,v1,y)
INTEGER m,n
REAL c2,dx,gamma,x1,y1,v1(1),y(3)
COMMON /sphcom/ c2,gamma,dx,m,n
  Supplies starting values for integration at  $x = -1 + dx$ .
y(3)=v1(1)
if(mod(n-m,2).eq.0)then
  y1=gamma
else
  y1=-gamma
endif
y(2)=-y(3)-c2)*y1/(2*(m+1))
y(1)=y1+y(2)*dx
return
END

SUBROUTINE load2(x2,v2,y)
INTEGER m,n
REAL c2,dx,gamma,x2,v2(2),y(3)
COMMON /sphcom/ c2,gamma,dx,m,n
  Supplies starting values for integration at  $x = 1 - dx$ .
y(3)=v2(2)
y(1)=v2(1)
y(2)=(y(3)-c2)*y(1)/(2*(m+1))
return
END

SUBROUTINE score(xf,y,f)
INTEGER i,m,n
REAL c2,gamma,dx,xf,f(3),y(3)
COMMON /sphcom/ c2,gamma,dx,m,n
  Tests whether solutions match at fitting point  $x = 0$ .
do 12 i=1,3
  f(i)=y(i)
enddo 12
return
END

```

## CITED REFERENCES AND FURTHER READING:

- Flammer, C. 1957, *Spheroidal Wave Functions* (Stanford, CA: Stanford University Press). [1]  
Abramowitz, M., and Stegun, I.A. 1964, *Handbook of Mathematical Functions*, Applied Mathematics Series, Volume 55 (Washington: National Bureau of Standards; reprinted 1968 by Dover Publications, New York), §21. [2]  
Morse, P.M., and Feshbach, H. 1953, *Methods of Theoretical Physics*, Part II (New York: McGraw-Hill), pp. 1502ff. [3]

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## 17.5 Automated Allocation of Mesh Points

In relaxation problems, you have to choose values for the independent variable at the mesh points. This is called *allocating* the grid or mesh. The usual procedure is to pick a plausible set of values and, if it works, to be content. If it doesn't work, increasing the number of points usually cures the problem.

If we know ahead of time where our solutions will be rapidly varying, we can put more grid points there and less elsewhere. Alternatively, we can solve the problem first on a uniform mesh and then examine the solution to see where we should add more points. We then repeat the solution with the improved grid. The object of the exercise is to allocate points in such a way as to represent the solution accurately.

It is also possible to automate the allocation of mesh points, so that it is done "dynamically" during the relaxation process. This powerful technique not only improves the accuracy of the relaxation method, but also (as we will see in the next section) allows internal singularities to be handled in quite a neat way. Here we learn how to accomplish the automatic allocation.

We want to focus attention on the independent variable  $x$ , and consider two alternative reparametrizations of it. The first, we term  $q$ ; this is just the coordinate corresponding to the mesh points themselves, so that  $q = 1$  at  $k = 1$ ,  $q = 2$  at  $k = 2$ , and so on. Between any two mesh points we have  $\Delta q = 1$ . In the change of independent variable in the ODEs from  $x$  to  $q$ ,

$$\frac{dy}{dx} = \mathbf{g} \quad (17.5.1)$$

becomes

$$\frac{dy}{dq} = \mathbf{g} \frac{dx}{dq} \quad (17.5.2)$$

In terms of  $q$ , equation (17.5.2) as an FDE might be written

$$\mathbf{y}_k - \mathbf{y}_{k-1} - \frac{1}{2} \left[ \left( \mathbf{g} \frac{dx}{dq} \right)_k + \left( \mathbf{g} \frac{dx}{dq} \right)_{k-1} \right] = 0 \quad (17.5.3)$$

or some related version. Note that  $dx/dq$  should accompany  $\mathbf{g}$ . The transformation between  $x$  and  $q$  depends only on the *Jacobian*  $dx/dq$ . Its reciprocal  $dq/dx$  is proportional to the density of mesh points.

Now, given the function  $\mathbf{y}(x)$ , or its approximation at the current stage of relaxation, we are supposed to have some idea of how we want to specify the density of mesh points. For example, we might want  $dq/dx$  to be larger where  $\mathbf{y}$  is changing rapidly, or near to the boundaries, or both. In fact, we can probably make up a formula for what we would like  $dq/dx$  to be proportional to. The problem is that we do not know the proportionality constant. That is, the formula that we might invent would not have the correct integral over the whole range of  $x$  so as to make  $q$  vary from 1 to  $M$ , according to its definition. To solve this problem we introduce a second reparametrization  $Q(q)$ , where  $Q$  is a new independent variable. The relation between  $Q$  and  $q$  is taken to be *linear*, so that a mesh spacing formula for  $dQ/dx$  differs only in its unknown proportionality constant. A linear relation implies

$$\frac{d^2 Q}{dq^2} = 0 \quad (17.5.4)$$

or, expressed in the usual manner as coupled first-order equations,

$$\frac{dQ(x)}{dq} = \psi \quad \frac{d\psi}{dq} = 0 \quad (17.5.5)$$

where  $\psi$  is a new intermediate variable. We add these two equations to the set of ODEs being solved.

Completing the prescription, we add a third ODE that is just our desired mesh-density function, namely

$$\phi(x) = \frac{dQ}{dx} = \frac{dQ}{dq} \frac{dq}{dx} \quad (17.5.6)$$

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where  $\phi(x)$  is chosen by us. Written in terms of the mesh variable  $q$ , this equation is

$$\frac{dx}{dq} = \frac{\psi}{\phi(x)} \quad (17.5.7)$$

Notice that  $\phi(x)$  should be chosen to be positive definite, so that the density of mesh points is everywhere positive. Otherwise (17.5.7) can have a zero in its denominator.

To use automated mesh spacing, you add the three ODEs (17.5.5) and (17.5.7) to your set of equations, i.e., to the array  $\mathbf{y}(j, k)$ . Now  $x$  becomes a dependent variable!  $Q$  and  $\psi$  also become new dependent variables. Normally, evaluating  $\phi$  requires little extra work since it will be composed from pieces of the  $g$ 's that exist anyway. The automated procedure allows one to investigate quickly how the numerical results might be affected by various strategies for mesh spacing. (A special case occurs if the desired mesh spacing function  $Q$  can be found analytically, i.e.,  $dQ/dx$  is directly integrable. Then, you need to add only two equations, those in 17.5.5, and two new variables  $x, \psi$ .)

As an example of a typical strategy for implementing this scheme, consider a system with one dependent variable  $y(x)$ . We could set

$$dQ = \frac{dx}{\Delta} + \frac{|d \ln y|}{\delta} \quad (17.5.8)$$

or

$$\phi(x) = \frac{dQ}{dx} = \frac{1}{\Delta} + \left| \frac{dy/dx}{y\delta} \right| \quad (17.5.9)$$

where  $\Delta$  and  $\delta$  are constants that we choose. The first term would give a uniform spacing in  $x$  if it alone were present. The second term forces more grid points to be used where  $y$  is changing rapidly. The constants act to make every logarithmic change in  $y$  of an amount  $\delta$  about as “attractive” to a grid point as a change in  $x$  of amount  $\Delta$ . You adjust the constants according to taste. Other strategies are possible, such as a logarithmic spacing in  $x$ , replacing  $dx$  in the first term with  $d \ln x$ .

#### CITED REFERENCES AND FURTHER READING:

Eggleton, P. P. 1971, *Monthly Notices of the Royal Astronomical Society*, vol. 151, pp. 351–364.  
 Kippenhan, R., Weigert, A., and Hofmeister, E. 1968, in *Methods in Computational Physics*, vol. 7 (New York: Academic Press), pp. 129ff.

## 17.6 Handling Internal Boundary Conditions or Singular Points

Singularities can occur in the interiors of two point boundary value problems. Typically, there is a point  $x_s$  at which a derivative must be evaluated by an expression of the form

$$S(x_s) = \frac{N(x_s, \mathbf{y})}{D(x_s, \mathbf{y})} \quad (17.6.1)$$

where the denominator  $D(x_s, \mathbf{y}) = 0$ . In physical problems with finite answers, singular points usually come with their own cure: Where  $D \rightarrow 0$ , there the physical solution  $\mathbf{y}$  must be such as to make  $N \rightarrow 0$  simultaneously, in such a way that the ratio takes on a meaningful value. This constraint on the solution  $\mathbf{y}$  is often called a *regularity condition*. The condition that  $D(x_s, \mathbf{y})$  satisfy some special constraint at  $x_s$  is entirely analogous to an extra boundary condition, an algebraic relation among the dependent variables that must hold at a point.

We discussed a related situation earlier, in §17.2, when we described the “fitting point method” to handle the task of integrating equations with singular behavior at the boundaries. In those problems you are unable to integrate from one side of the domain to the other.

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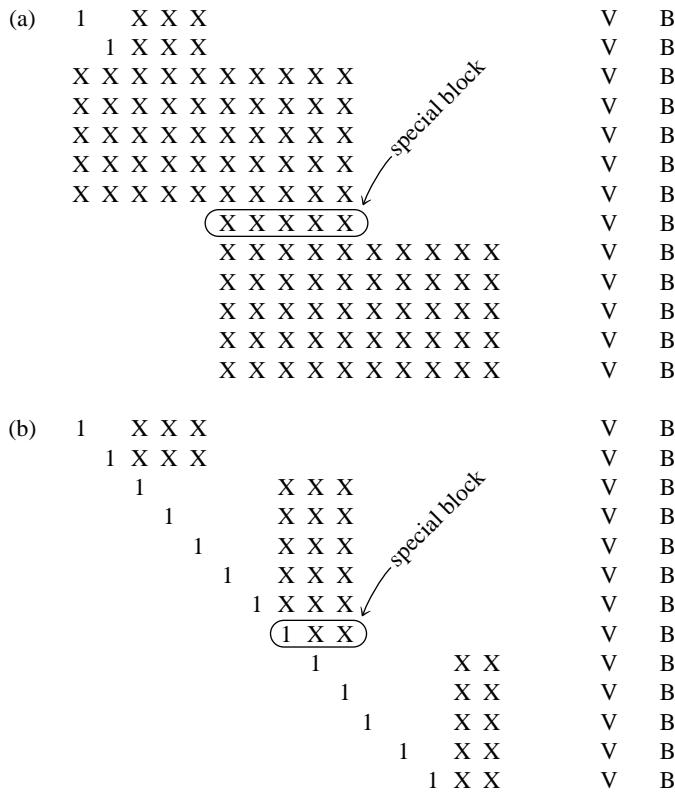


Figure 17.6.1. FDE matrix structure with an internal boundary condition. The internal condition introduces a special block. (a) Original form, compare with Figure 17.3.1; (b) final form, compare with Figure 17.3.2.

However, the ODEs do have well-behaved derivatives and solutions in the neighborhood of the singularity, so it is readily possible to integrate away from the point. Both the relaxation method and the method of “shooting” to a fitting point handle such problems easily. Also, in those problems the presence of singular behavior served to isolate some special boundary values that had to be satisfied to solve the equations.

The difference here is that we are concerned with singularities arising at intermediate points, where the location of the singular point depends on the solution, so is not known *a priori*. Consequently, we face a circular task: The singularity prevents us from finding a numerical solution, but we need a numerical solution to find its location. Such singularities are also associated with selecting a special value for some variable which allows the solution to satisfy the regularity condition at the singular point. Thus, internal singularities take on aspects of being internal boundary conditions.

One way of handling internal singularities is to treat the problem as a free boundary problem, as discussed at the end of §17.0. Suppose, as a simple example, we consider the equation

$$\frac{dy}{dx} = \frac{N(x, y)}{D(x, y)} \tag{17.6.2}$$

where  $N$  and  $D$  are required to pass through zero at some unknown point  $x_s$ . We add the equation

$$z \equiv x_s - x_1 \quad \frac{dz}{dx} = 0 \tag{17.6.3}$$

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where  $x_s$  is the unknown location of the singularity, and change the independent variable to  $t$  by setting

$$x - x_1 = tz, \quad 0 \leq t \leq 1 \quad (17.6.4)$$

The boundary conditions at  $t = 1$  become

$$N(x, y) = 0, \quad D(x, y) = 0 \quad (17.6.5)$$

Use of an adaptive mesh as discussed in the previous section is another way to overcome the difficulties of an internal singularity. For the problem (17.6.2), we add the mesh spacing equations

$$\frac{dQ}{dq} = \psi \quad (17.6.6)$$

$$\frac{d\psi}{dq} = 0 \quad (17.6.7)$$

with a simple mesh spacing function that maps  $x$  uniformly into  $q$ , where  $q$  runs from 1 to  $M$ , the number of mesh points:

$$Q(x) = x - x_1, \quad \frac{dQ}{dx} = 1 \quad (17.6.8)$$

Having added three first-order differential equations, we must also add their corresponding boundary conditions. If there were no singularity, these could simply be

$$\text{at } q = 1 : \quad x = x_1, \quad Q = 0 \quad (17.6.9)$$

$$\text{at } q = M : \quad x = x_2 \quad (17.6.10)$$

and a total of  $N$  values  $y_i$  specified at  $q = 1$ . In this case the problem is essentially an initial value problem with all boundary conditions specified at  $x_1$  and the mesh spacing function is superfluous.

However, in the actual case at hand we impose the conditions

$$\text{at } q = 1 : \quad x = x_1, \quad Q = 0 \quad (17.6.11)$$

$$\text{at } q = M : \quad N(x, y) = 0, \quad D(x, y) = 0 \quad (17.6.12)$$

and  $N - 1$  values  $y_i$  at  $q = 1$ . The “missing”  $y_i$  is to be adjusted, in other words, so as to make the solution go through the singular point in a regular (zero-over-zero) rather than irregular (finite-over-zero) manner. Notice also that these boundary conditions do not directly impose a value for  $x_2$ , which becomes an adjustable parameter that the code varies in an attempt to match the regularity condition.

In this example the singularity occurred at a boundary, and the complication arose because the location of the boundary was unknown. In other problems we might wish to continue the integration beyond the internal singularity. For the example given above, we could simply integrate the ODEs to the singular point, then as a separate problem recommence the integration from the singular point on as far we care to go. However, in other cases the singularity occurs internally, but does not completely determine the problem: There are still some more boundary conditions to be satisfied further along in the mesh. Such cases present no difficulty in principle, but do require some adaptation of the relaxation code given in §17.3. In effect all you need to do is to add a “special” block of equations at the mesh point where the internal boundary conditions occur, and do the proper bookkeeping.

Figure 17.6.1 illustrates a concrete example where the overall problem contains 5 equations with 2 boundary conditions at the first point, one “internal” boundary condition, and two final boundary conditions. The figure shows the structure of the overall matrix equations along the diagonal in the vicinity of the special block. In the middle of the domain, blocks typically involve 5 equations (rows) in 10 unknowns (columns). For each block prior to the special block, the initial boundary conditions provided enough information to zero the first two columns of the blocks. The five FDEs eliminate five more columns, and the final three columns need to be stored for the backsubstitution step (as described in §17.3). To handle the extra condition we break the normal cycle and add a special block with only one equation:

the internal boundary condition. This effectively reduces the required storage of unreduced coefficients by one column for the rest of the grid, and allows us to reduce to zero the first three columns of subsequent blocks. The subroutines `red`, `pinvs`, `bksub` can readily handle these cases with minor recoding, but each problem makes for a special case, and you will have to make the modifications as required.

#### CITED REFERENCES AND FURTHER READING:

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